

Courant Institute of  
Mathematical Sciences

AEC Computing and Applied Mathematics Center

A 16 Channel Medium Speed Multiplexer

Vincent T. Prestianni

AEC Research and Development Report

Mathematics and Computers  
May 1971



New York University

NYO-1480-171 C.2



UNCLASSIFIED

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Contract No. AT(30-1)-1480

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## 1. INTRODUCTION

A medium speed multiplexer system that interfaces 16 synchronous serial channels to a Honeywell 16-bit series computer is described in this report. The system permits simultaneous communication between the Control Data 6600 computer, at the Courant Institute of Mathematical Sciences, and any 16 sites accessible by telephone or private line.

The salient features of the system are:

- a. The multiplexer channels function with any synchronous data device in a half or full duplex mode.
- b. The multiplexer channels handle data rates up to 10k bps.
- c. The multiplexer channels are individually synchronized by the data channel clocks.
- d. Eight-bit byte registration for a channel in the input mode is accomplished within the multiplexer. This feature can be bypassed for any channel if complete data transparency is desired.
- e. Software overhead is less for 16 channel communication with the multiplexer/computer composite than with the computer alone and attendant interfaces.
- f. The operating modes of each channel are controlled by the Honeywell computer software.

g. The data passing through the multiplexer, whether transmitting or receiving on any channel, is in no way altered. The philosophy of data transparency, a necessity in a research environment, is thus adhered to.

The multiplexer handles all remote batch entry of the CDC-6600. Figure 1 shows the CDC-6600/DDP-516 configuration\* at the Computing Center.

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\*The CDC-6600/DDP-516 interface and the remote entry system for the CDC-6600 are described in the following AEC Research and Development Reports published by New York University: Bianchini, R. P., "An Interface Between a Control Data 6600 Series Computer and a Honeywell 16-Bit Series Computer", NYO-1480-119, October 1969; Francheschini, E., Feinroth, Y., and Goldstein, M., "A Remote Entry System for the CDC-6600", NYO-1480-148, April 1970.



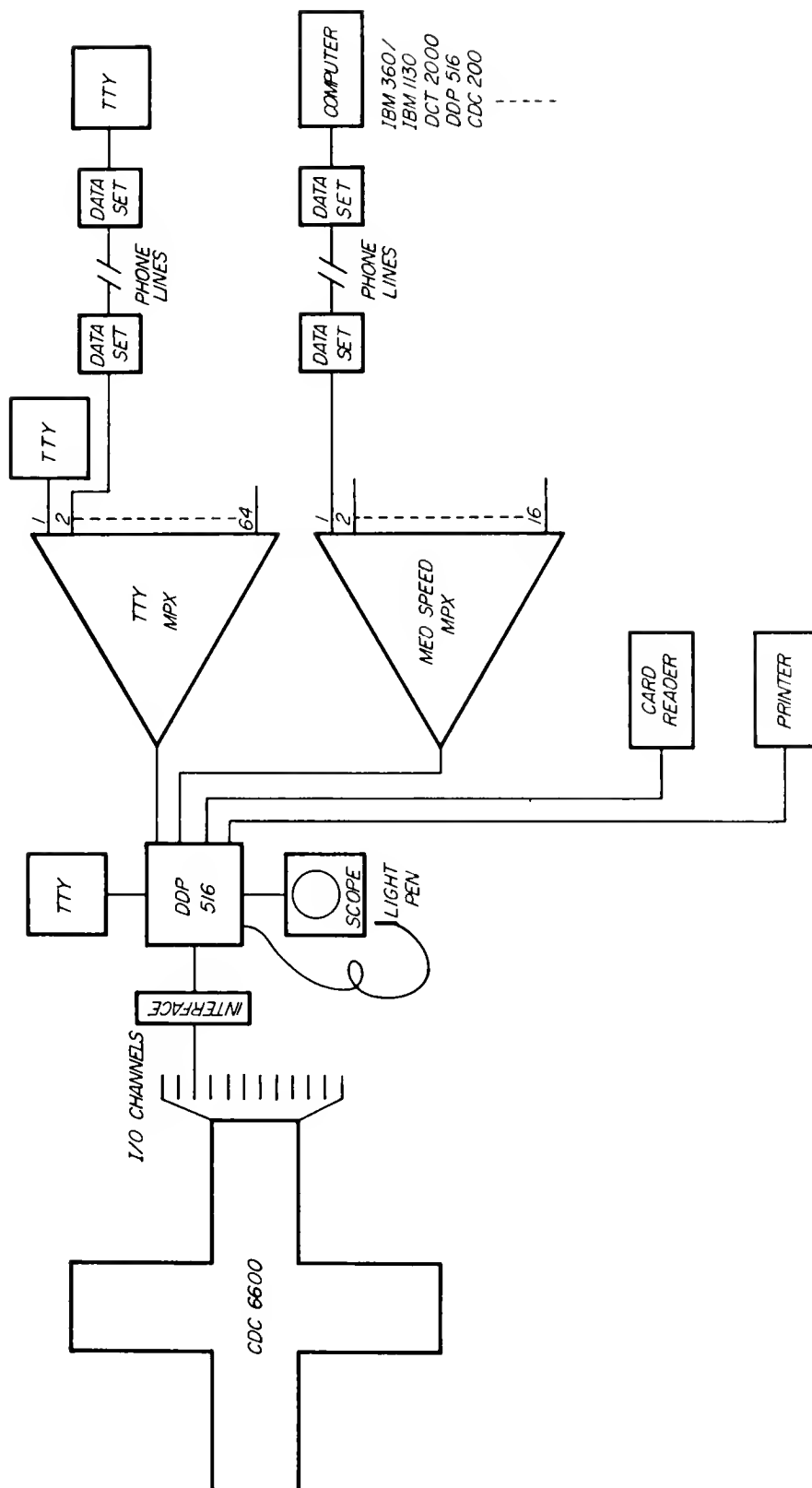


FIGURE 1 - 6600/516 CONFIGURATION

## 2. SYSTEM OPERATION

The remote entry configuration of the CDC-6600 consists of 16 synchronous serial channels interfaced by the multiplexer to a Honeywell 16-bit computer. The Honeywell computer is interfaced to the CDC-6600. This small "front end" computer, which is interfaced to the CDC-6600, performs the dual function of I/O controller and temporary storage. Any of the Honeywell models DDP-116/416/516 or H-316 can be used for this purpose.

The data to and from the serial channels is in the form of bit streams driven by channel clocks. The I/O structure of the Honeywell computer permits communication on a byte level. Therefore, the throughput of the multiplexer requires that it contain a memory. Since the multiplexer/computer communication is asynchronous, the multiplexer must be capable of communicating to the computer the real time requirements of the 16 serial channels. Accordingly, the multiplexer utilizes the hardware interrupt feature of the computer.

From an overall system point of view, the computer outputs data through the multiplexer to any one of the 16 serial channels. Transmission from computer to multiplexer is in 8-bit bytes. The multiplexer distributes the bytes as specified by the computer. Then the 8-bit bytes are marched out serially onto the appropriate channels. Since the channels are synchronized to channel trans-

mission clocks, the computer must maintain an output data rate of one 8-bit byte for every 8 counts of the corresponding channel transmission clock.

The computer collects data from any one of the 16 serial channels. The bit streams are synchronized to channel derived receive clocks. The multiplexer collects and stores the bits into groups of 8-bit bytes. The computer takes these bytes along with additional information for channel identification. Since the channels are synchronized to channel derived receive clocks, the computer must maintain an input data rate of one 8-bit byte for every 8 counts of the corresponding channel clock. While the bits are collected, the multiplexer hardware establishes the byte synchronization of the received message. The computer software is thus relieved of this task. A hardware switch can bypass this feature for any channel.

The multiplexer allows each of the 16 channels to operate independently. Thus, each channel may operate in a receive or transmit mode, at different data rates, in full or half-duplex, and with or without byte synchronization for received messages.

The computer interrupt circuitry controls the byte level communication between the computer and multiplexer. When the computer is interrupted for input, it collects the 8-bit bytes that the receive channels have ready. When the computer is interrupted for output, it outputs the 8-bit bytes that the multiplexer can accept for the channels in the transmit mode.

The multiplexer memory contains 48 bits for each channel: 24 for input and 24 for output data. In each case, double level byte buffering is employed. The design allows the computer an 8 bit time interval for servicing the multiplexer for input or output data. A real time hardware clock in the multiplexer paces the computer input and output interrupts. The real time bit requirements of the 16 serial channels determine the pacing rate.

Previous techniques would have required the computer to treat the 16 serial channels with individual 8-bit byte interfaces as separate devices, all with separate interrupts. This creates a software environment with a high interrupt overhead. The multiplexer has only two interrupts for all of the 16 serial channels: one for input, one for output. This feature reduces the interrupt software overhead that is required for communication between the 16 serial channels and the computer.

### 3. HARDWARE STRUCTURE

The multiplexer comprises four sections--see Fig. 2.

A. Clock, Timing, and Decode Circuitry. Synchronizes the delay line memory to the pre-multiplexer scanning and to the computer/multiplexer interface.

B. Pre-Multiplexer. Synchronizes each serial data channel to the multiplexer memory. Contains the individual channel mode selection switches.

C. Delay Line Memory. Consists of twelve 1 MHz delay lines for double level buffering of each of 16 channels. Each delay line has a 64 bit capacity.

D. Computer/Multiplexer Interface. Functions as a communication channel between the delay line memory and the computer. Can buffer an 8-bit byte for computer input and output data. Buffering is necessary because of the 64  $\mu$ sec access time of the delay line memory.

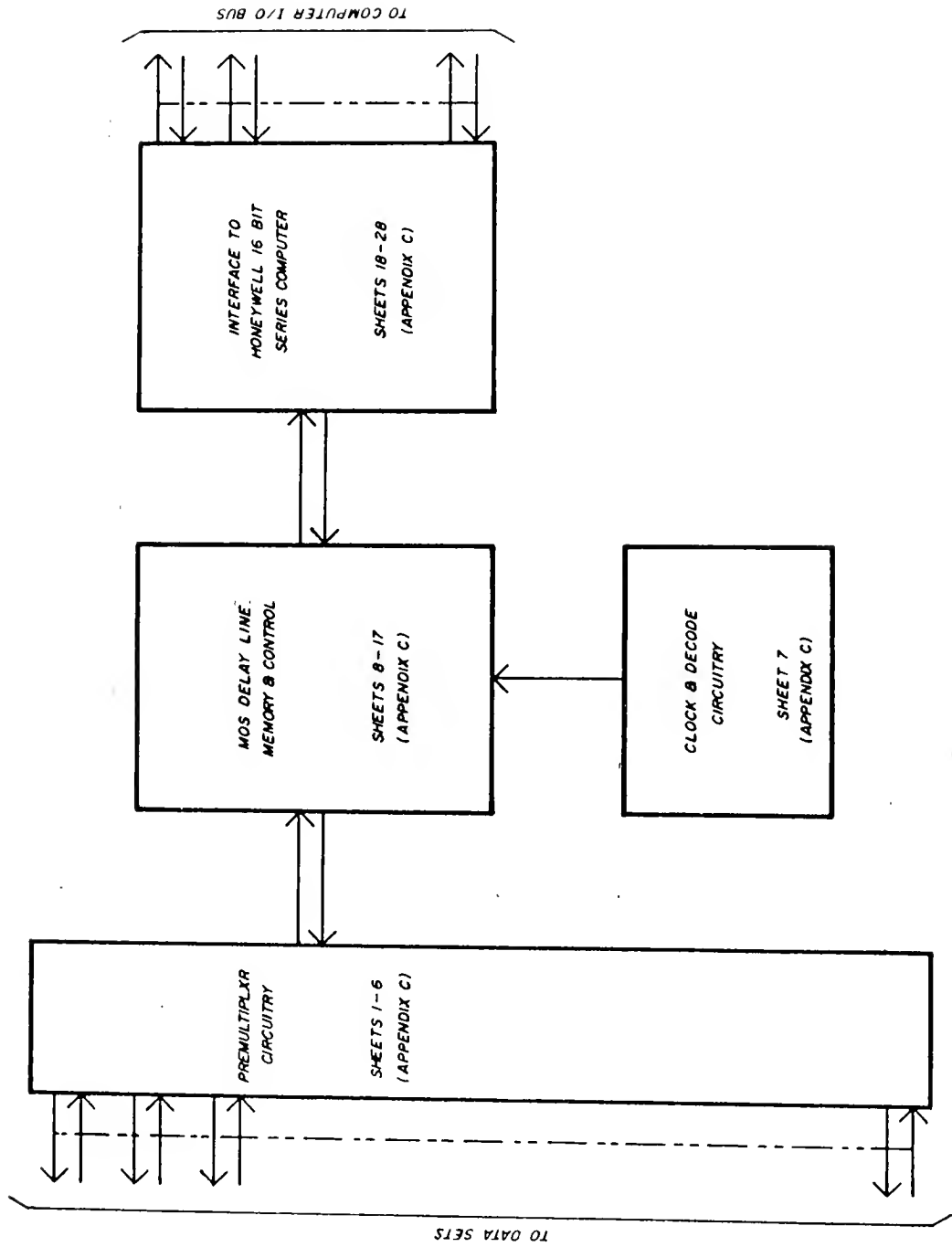


FIGURE 2 - SYSTEM BLOCK DIAGRAM OF 16 CHANNEL MULTIPLEXER

### 3A. Clock, Timing, and Decode Circuitry

(See Appendix C Schematics, Sheet no. 7)

Contains a 1 MHz crystal control clock, a 6 bit countdown chain, a 2 bit decoder, and a 4 bit decoder. The 2 bit decoder cycles every 4 clock counts; the 4 bit decoder cycles every 64 clock counts. See Fig. 3. Since the 1 MHz clock drives the delay lines, one cycle of the 4 bit decoder corresponds to one delay line trip. Fig. 4 shows the timing relationship of all the clock outputs. The 4 bit decoder has sixteen 4  $\mu$ sec states: 00<sub>8</sub>, 01<sub>8</sub>, 02<sub>8</sub>, 03<sub>8</sub>, ..., 17<sub>8</sub>. Each state has four 1  $\mu$ sec sub-states: T0, T1, T2, and T3. Thus in the sequential time allocation of the 16 data set channels, the time slice for each is subdivided into four 1  $\mu$ sec sub-states: T0, T1, T2, and T3.

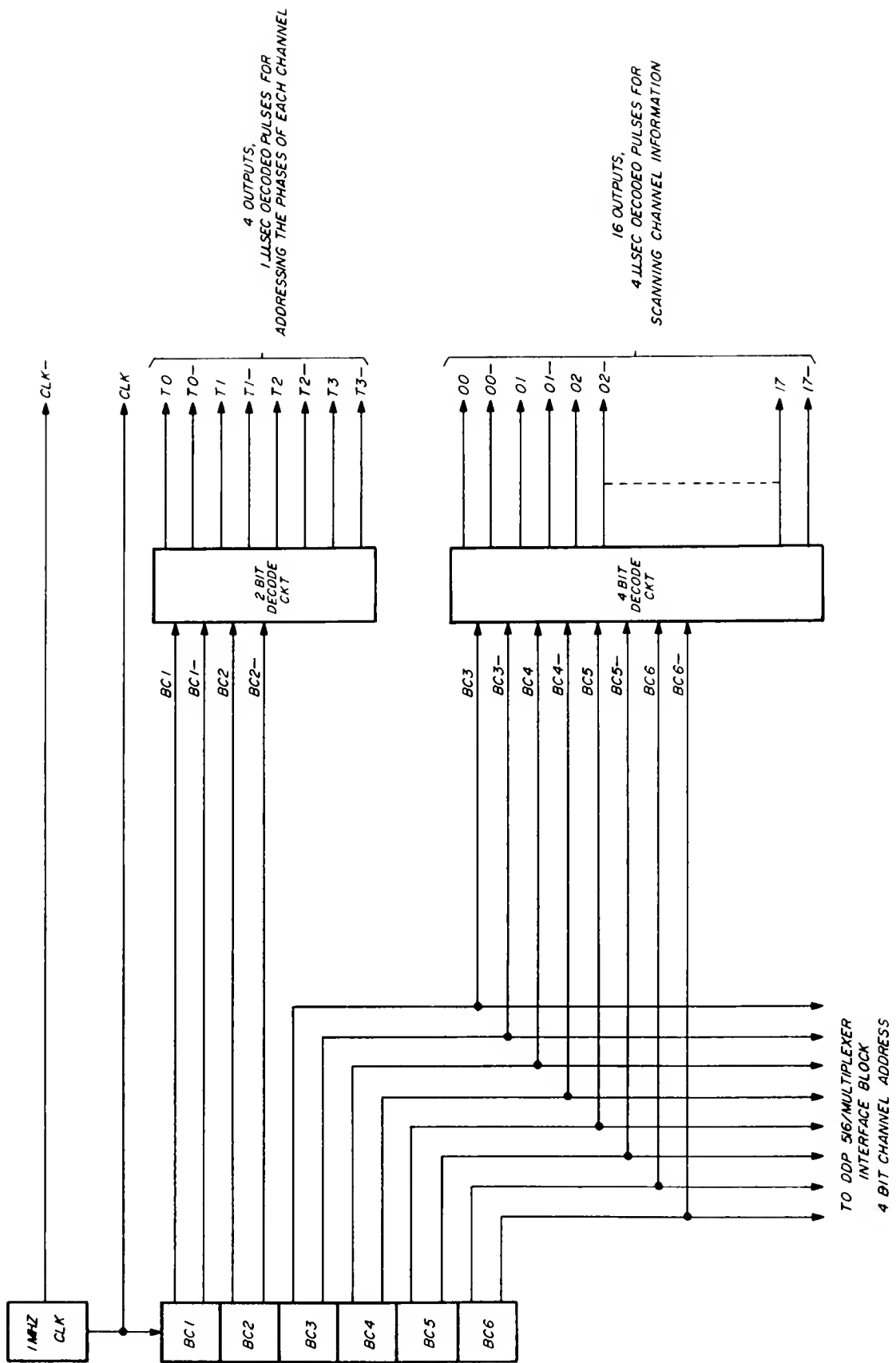


FIGURE 3 - SYSTEM CLOCK AND DECODE CIRCUITRY



# TIMING WAVESHAPES GENERATED BY CLK, DEC CIRCUITRY

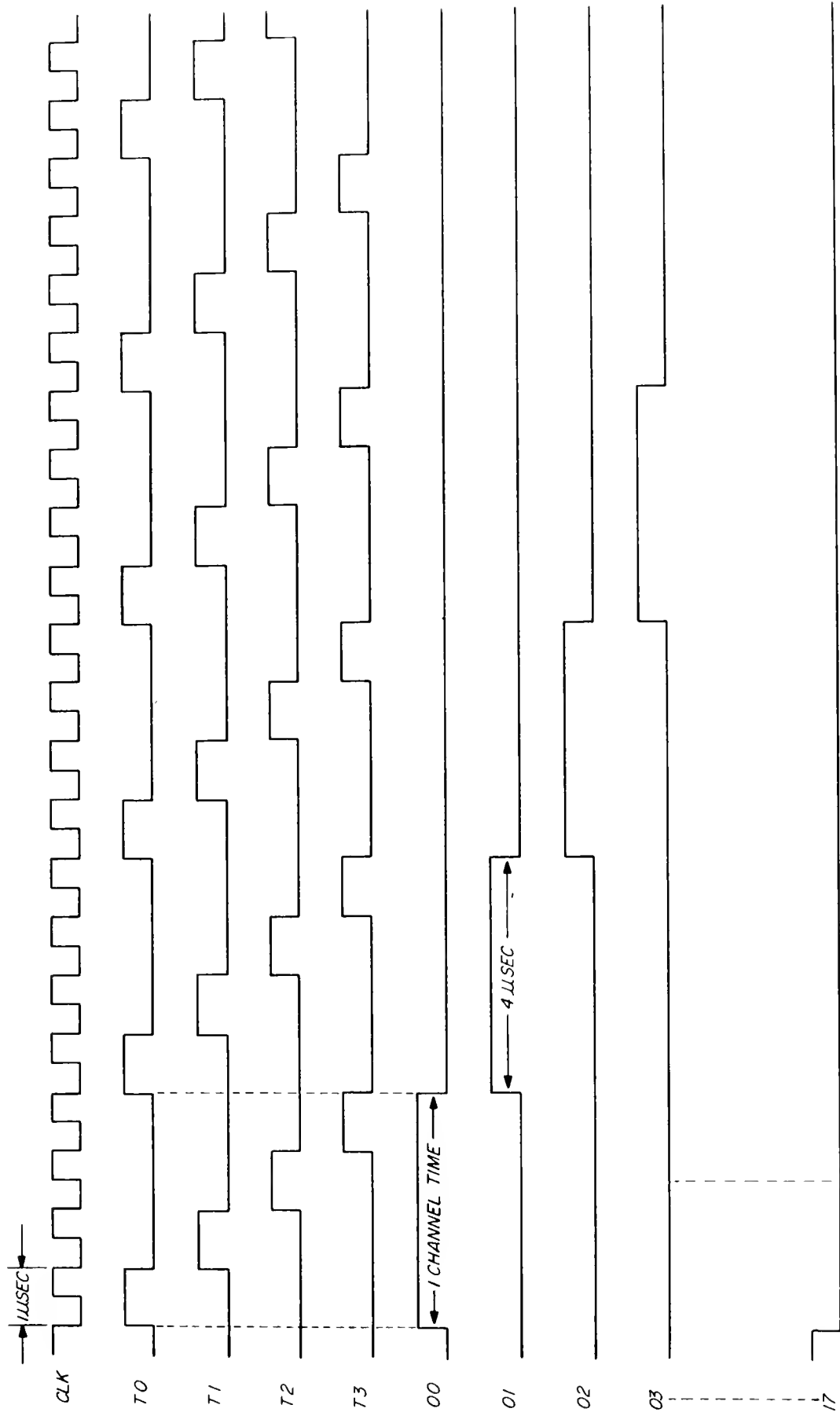


FIGURE 4

### 3B. Pre-Multiplexer Circuitry

(Appendix C Schematics, Sheets 1-6)

Each pre-multiplexer channel comprises seven flip-flops, "nand" gates for scanning and control, and level convertor circuits for voltage compatability between the pre-multiplexer and the synchronous serial data channels.

Two flip-flops are used for channel mode selection, and are separately controlled by the computer software. Two flip-flops are used for receiving serial data from the channel. Three flip-flops are used for transmitting serial data onto the channel.

#### Data Terminal Ready Flip-Flop

In the "1" state, enables the channel to operate. In the "0" state, closes off the logic and disconnects the channel from the remote user.

#### Request to Send Flip-Flop

In the "1" state, sets the channel into the transmit mode. In the "0" state, sets the channel into the receive mode. The state of this flip-flop is forced by a hardware switch when the channel is used in a full-duplex mode.

### Receive Data and Receive Clock Flip-Flops

The receive data flip-flop functions as a one-bit serial buffer. The receive clock flip-flop functions as a switch which is set by the channel receive clock. The switch remains set until the delay line memory can accept the bit from the receive data flip-flop. This can occur only during the 4  $\mu$ sec time slice assigned to the channel.

### Transmit Clock and Two Transmit Data Flip-Flops

The transmit clock flip-flop functions as a switch which is set by the channel transmit clock. The switch remains set until the delay line memory can output data to the two transmit data flip-flops. The two flip-flops function as a two bit serial buffer which outputs data to the channel at a rate determined by the transmit clock.

### 3C. Delay Line Memory

(Appendix C Schematics, Sheets 8-17)

This section of the hardware contains the memory required for control and double level buffering of all of the 16 serial data channels. The receive buffers are used with a compare circuit to establish byte synchronization of the receive message. The memory consists of 64 sequential 12-bit words. A new word is accessed from the sequential memory at a 1- $\mu$ sec rate. The total cycling time of the memory is 64  $\mu$ seconds. Four of the 12-bit words are used for handling data bits of each serial data channel.

Channel bit storage requirements are shown in Tables I and II, p. 17-19. Delay line memory bit assignments are shown in Fig. 5. The first two words, T0 and T1, are used for channel receive data. The second two words, T2 and T3, are used for channel transmit data. Because of the sequential nature of the memory, byte transfers from T0 to T1 and T2 to T3 are accomplished by using an 8-bit data delay register (8 flip-flops). Accordingly, in the receive mode, the channel data collection buffer T0 precedes the byte static buffer T1; in the transmit mode, the channel data distribution buffer T3 follows the static buffer T2.

The combination of serial and parallel 8-bit buffers, for input and output, forms the double level buffering mechanism for each channel. The byte static buffers store 8-bit bytes. The serial

buffers satisfy the real time requirements of the serial data channels (one bit per channel clock interval).

The delay line memory consists of the following circuitry plus interconnecting logic and flip-flops--see Fig. 6.

### 8-Bit Data Register

This register is inserted between the output and input of the delay line memory. It is used to transfer byte data from the serial buffer to the byte static buffer (receive mode), or from the byte static buffer to the serial buffer (transmit mode). These transfers take place when the input serial buffer goes full, or when the output serial buffer goes empty.

### Synch Compare Circuitry

Establishes the first bit of a receive message. The circuit compares the contents of the receive serial buffer against a known pattern (synch), as the receive bit stream is marching through the receive serial buffer.

### Add "1" Logic

Six bits for each channel are reserved in memory to store two separate 3 bit counts. One set is for the number of receive serial buffer shifts, the other is for the number of transmit serial buffer shifts. The add "1" logic updates these counters every time data is shifted.

### Twelve 64-Bit Delay Lines

Each delay line is a 64-bit MOS 1 MHz shift register. A stored bit can only be read when it is shifted out of the end of the delay line. By feeding the delay line output back into the input logic gates, a bit can either be stored indefinitely or modified as required. An expandable "or" node is provided at the input of each delay line. This allows for the addition of as many control functions as needed. The 12 delay lines operate in parallel; therefore, at any one time, a 12-bit word is read out by looking at the output of each of the delay lines. Because of the output to input feedback, this word can be either stored indefinitely or modified. When clocking these lines at a 1 MHz rate, a new 12-bit word appears at the outputs of the delay lines every one  $\mu$ second. Since these lines are 64 bits long, a particular 12-bit word is accessible every 64  $\mu$ seconds. At this point, it can either be modified or reinserted as is into the lines. The lines operate in complete synchronism with the scanning of the 16 pre-multiplexer channels. Each pre-multiplexer is allowed a 4  $\mu$ sec time slice--Fig. 7. During this time slice 4 sequential, 1  $\mu$ sec 12-bit words of delay line memory are accessible. This occurs for each pre-multiplexer channel every 64  $\mu$ seconds.

Table I - Input Data Bits For Each Serial Channel

1. 7 bits are used as a serially loaded, parallel-out shift register buffer. The buffer shifts in data at a rate equal to the receive channel clock.
2. 8 bits are used as a parallel-in, parallel-out byte static buffer.
3. 3 bits are used to store the number of shifts of the receive serially loaded buffer. The count is updated by 1 for each shift of the serial buffer.
4. Full bit. This is a switch which is set when the serial buffer is full, and its contents are transferred to the static buffer. This occurs when the 3-bit stored count cycles through its zero state. The switch is reset when the multiplexer/computer interface can accept the 8-bit byte from the byte static buffer.
5. 2 bits are used to establish the first bit of the receive message. These are the pre-synch and synch bits.
6. 2 bits are used to mark the first bit of the receive message, and to send a termination word to the computer when the serial

channel is no longer receiving data. These are the start and stop bits.

7. Overwrite bit. This is a switch which is set when the computer fails to meet the receive real time byte requirements of the serial channel.



Table II - Output Data Bits For Each Serial Channel

1. 8 bits are used as a parallel-in, parallel-out byte buffer.
2. 8 bits are used as a parallel-in, serial output shift register. This buffer shifts out data at a rate equal to the transmit channel clock.
3. 3 bits are used to store the number of shifts of the transmit serial output buffer. The count is updated by 1 for each shift.
4. Empty bit. This is a switch which is set when the serial output buffer is empty. At this point, an 8-bit byte is transferred to it from the byte static buffer. This occurs when the 3-bit stored count cycles through its zero state. The switch is reset when the multiplexer/computer interface loads the static buffer with another 8-bit byte.
5. Underwrite bit. This is a switch which is set when the computer fails to meet the transmit real time byte requirement of the serial channel.

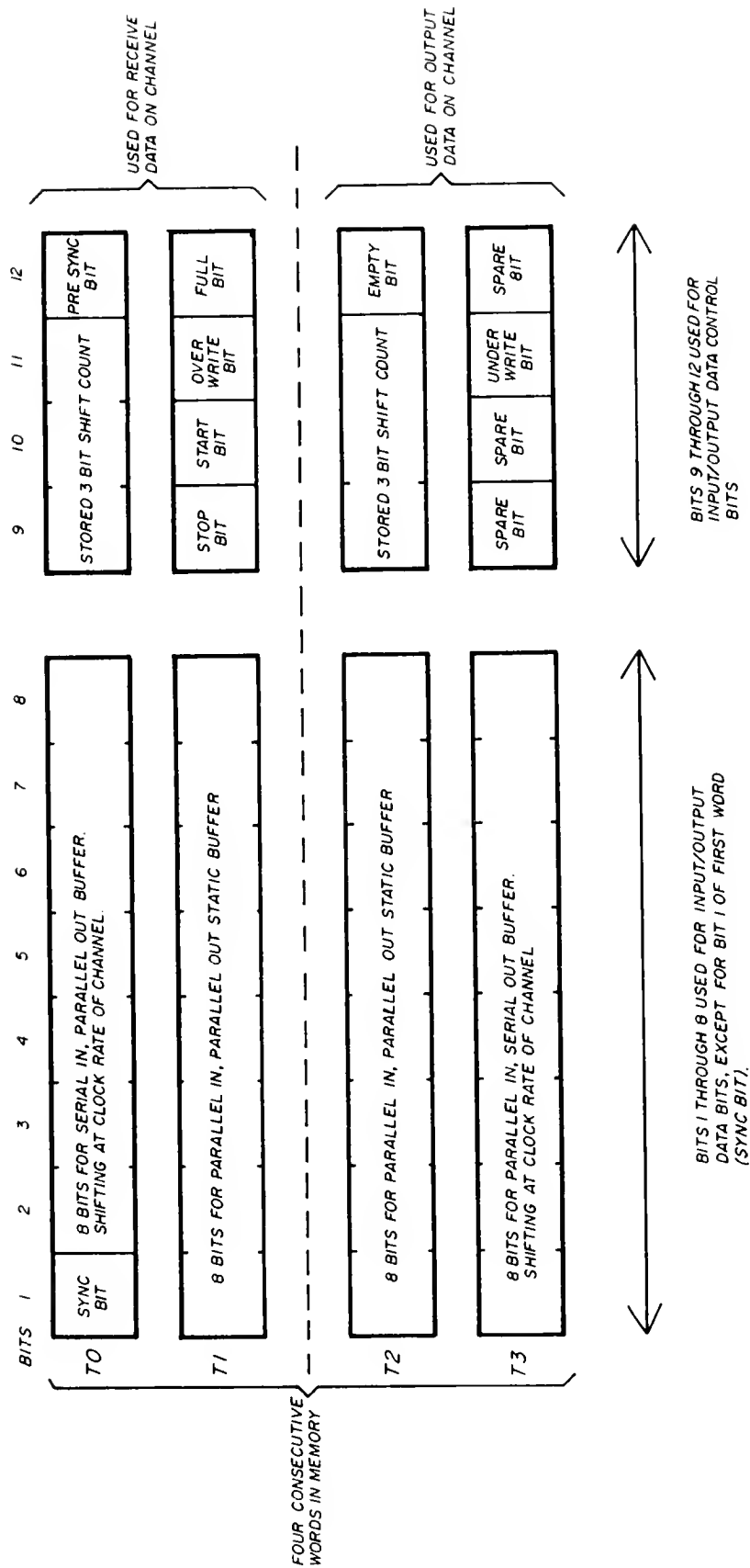


FIGURE 5 - BIT ASSIGNMENTS IN MEMORY FOR HANDLING ONE OF SIXTEEN SERIAL CHANNELS. TOTAL MEMORY REQUIREMENT IS 64 TWELVE BIT WORDS.

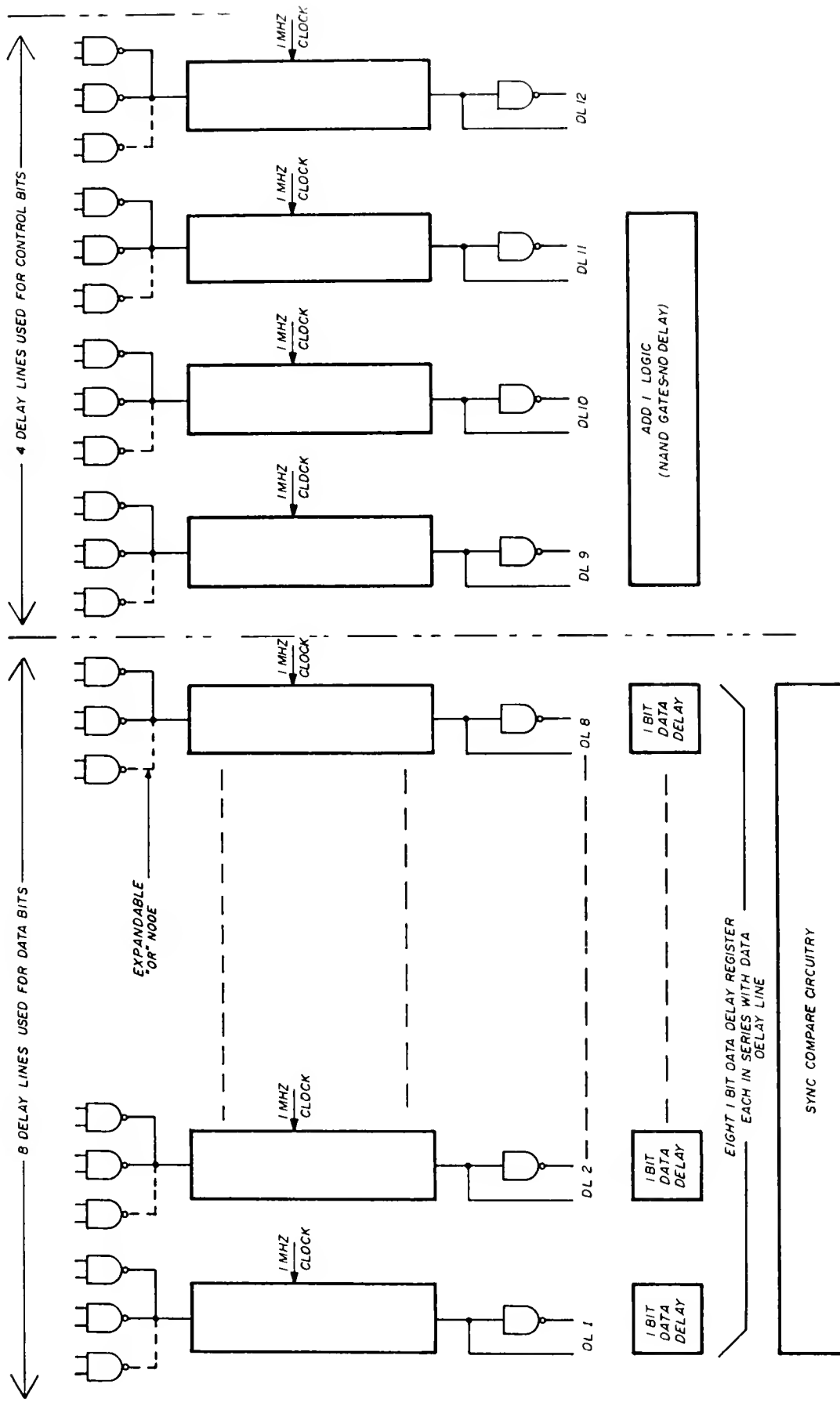
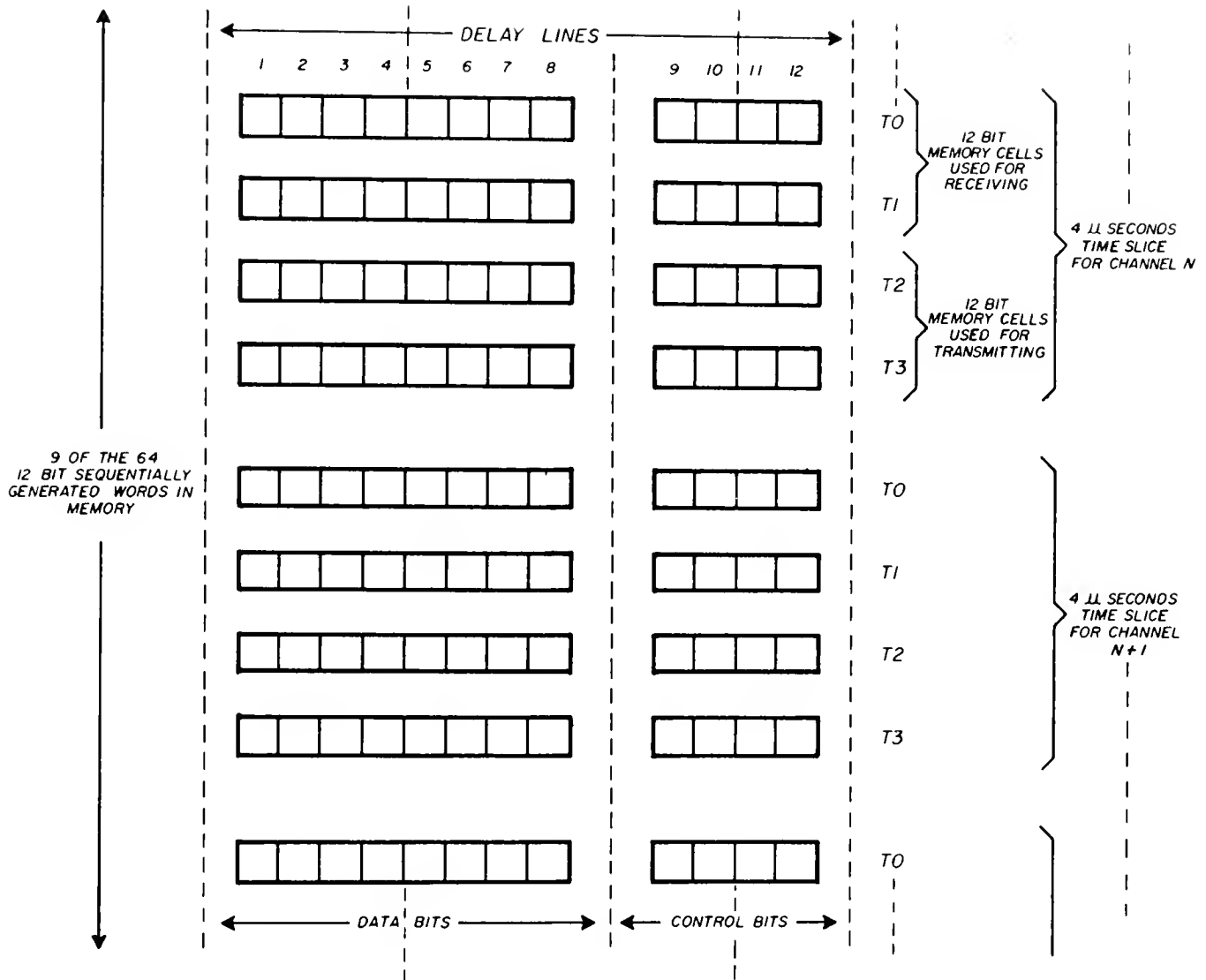


FIGURE 6 -- DELAY LINE MEMORY



MEMORY SYSTEM IS SYNCHRONIZED TO THE PRE-MULTIPLEXER CHANNEL SCANNING SYSTEM. THEREFORE A GIVEN PRE-MULTIPLEXER CHANNEL HAS ACCESS TO ONLY ONE GROUP OF FOUR  $\mu$ SEC WORDS (T0,T1,T2,T3). 16 CHANNELS ARE HANDLED BY THIS SEQUENTIALLY CONSTRUCTED MEMORY.

FIGURE 7 - MEMORY SYNCHRONIZATION

### 3D. Computer/Multiplexer Interface

(Appendix C Schematics, Sheets 18-28)

This section of the hardware contains two interfaces. They interface the delay line memory receive and transmit byte static buffers to the I/O bus of the computer. One interface contains the input data register, and the other the output data register. These registers provide the computer I/O bus with access to the delay line memory. A real time clock is used to pace the computer hardware interrupts for the 2 registers. The frequency of the clock depends upon the real time requirements of the channel with the highest bit rate.

#### Input Data Register System (Appendix C Schematics, Sheets 18, 19, 23)

When the byte storage buffers (word T1, Fig. 7) of the delay line memory become full, they compete for the input data register. When one byte is in this register, bytes are gated out until the register goes empty. The full condition of this register is gated with the pacing circuit. This generates an interrupt that causes the computer to enter an input routine.

When the input data register goes empty, it is ready to accept new bytes from the delay line memory. The process of the delay line memory emptying its byte storage buffers into the input data register continues until all receive byte buffers are empty.

As byte data is transferred from the delay line memory to the input data register, 4 bits from the clock, timing, and decode circuitry are strobed into 4 flip-flops of the register for channel identification. An additional 4 bits are strobed in the register. These are the overwrite, underwrite, start, and stop bits--see Fig. 5.

The overwrite and underwrite bits are set only when the computer fails to satisfy the real time byte requirements of the channel for either the receive or transmit mode. The start bit is set to indicate to the computer the first byte of the receive message. The stop bit is set when the channel stops receiving.

#### Output Data Register (Appendix C Schematics, Sheets 20, 25-27)

When the byte static buffers (word T2, Fig. 7) of the delay line memory go empty, an empty-flag is set in the 16-bit delay line status register--Fig. 9. The empty-flag corresponds to the channel that went empty. The flag "or" output is gated with the pacing circuit for computer interrupt operation. The computer senses this empty-flag, and the empty/full status of the output data register--Fig. 9. When the register is empty and the empty-flag is set, the computer is allowed to output a word for the corresponding transmitting channel. The computer output word is in the format shown in Fig. 9. The successful execution of this output resets the empty-flag in the 16-bit delay line status register. The empty-flag that gets reset is determined by the transmit channel 4-bit address, as indicated in the word format.

With the 16-bit delay line status register, the computer does not have to access the delay line memory (64  $\mu$ sec access time) to check the empty/full status of a given byte static buffer. The register allows the computer to sense the condition without any memory (delay lines) access time. The importance of knowing this condition in advance is to prevent the computer from placing data, which the delay line memory will not accept, in the output data register.

The output data register, which stores byte data with its corresponding 4-bit channel address, remains full until the appropriate cell in the delay line memory is addressed by using the address compare circuit--Fig. 9. Since the total cycle time of the memory is 64  $\mu$ seconds, the byte data transfer to memory will take place within this time. When the output data register goes empty, it is ready to accept new computer output data.

### Interrupt Operation

There are two separate interrupts generated: one for all channels with receive byte data, and the other for all channels with transmit byte data. A real time clock is used to pace the interrupts.

The computer can then service the multiplexer in one routine for all transmitting channels, and one other routine for all receiving channels. This reduces the software overhead required for handling 16 serial data channels. The frequency of the interrupt

pacers must be selected to satisfy the byte level requirements of the serial channel with the highest bit rate.



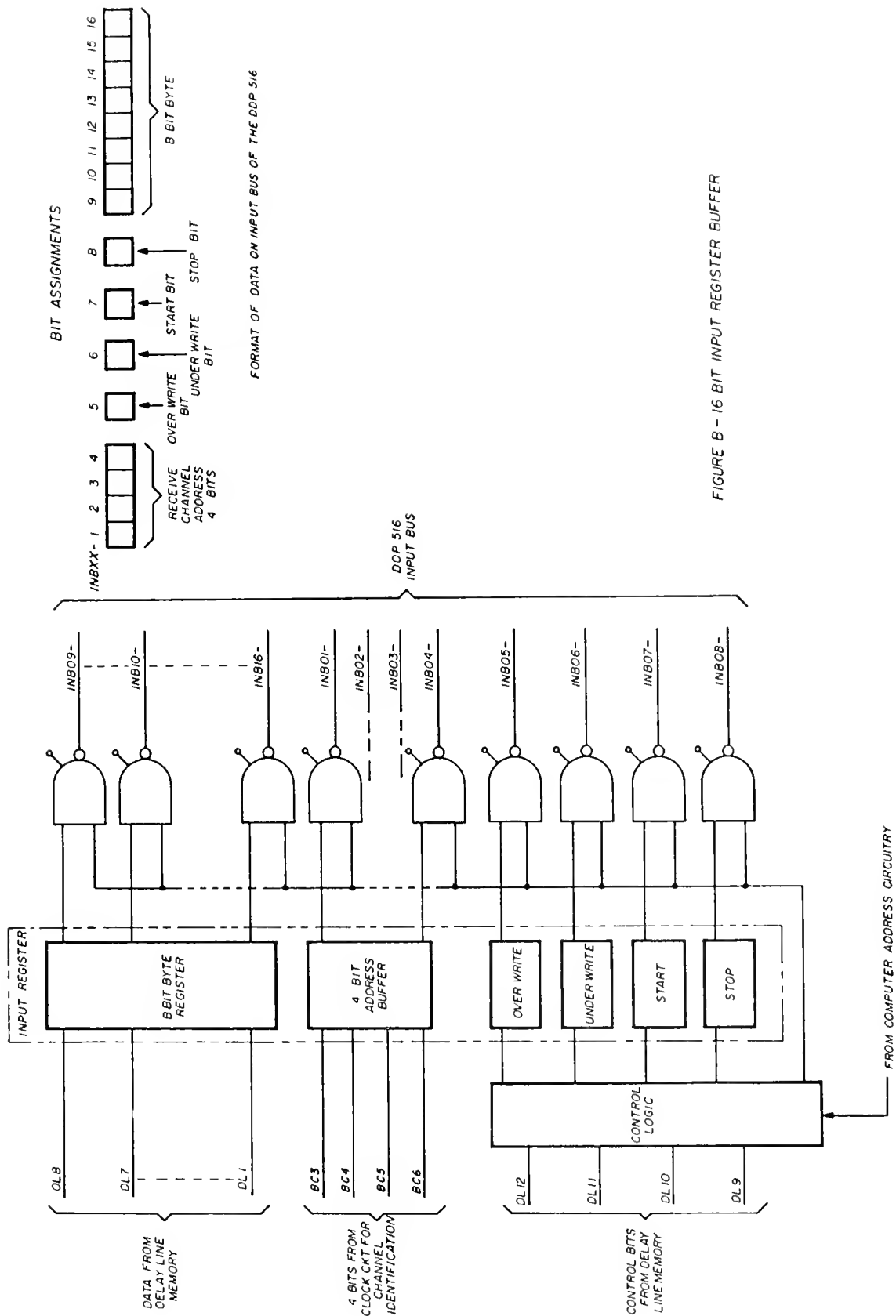


FIGURE B - 16 BIT INPUT REGISTER BUFFER

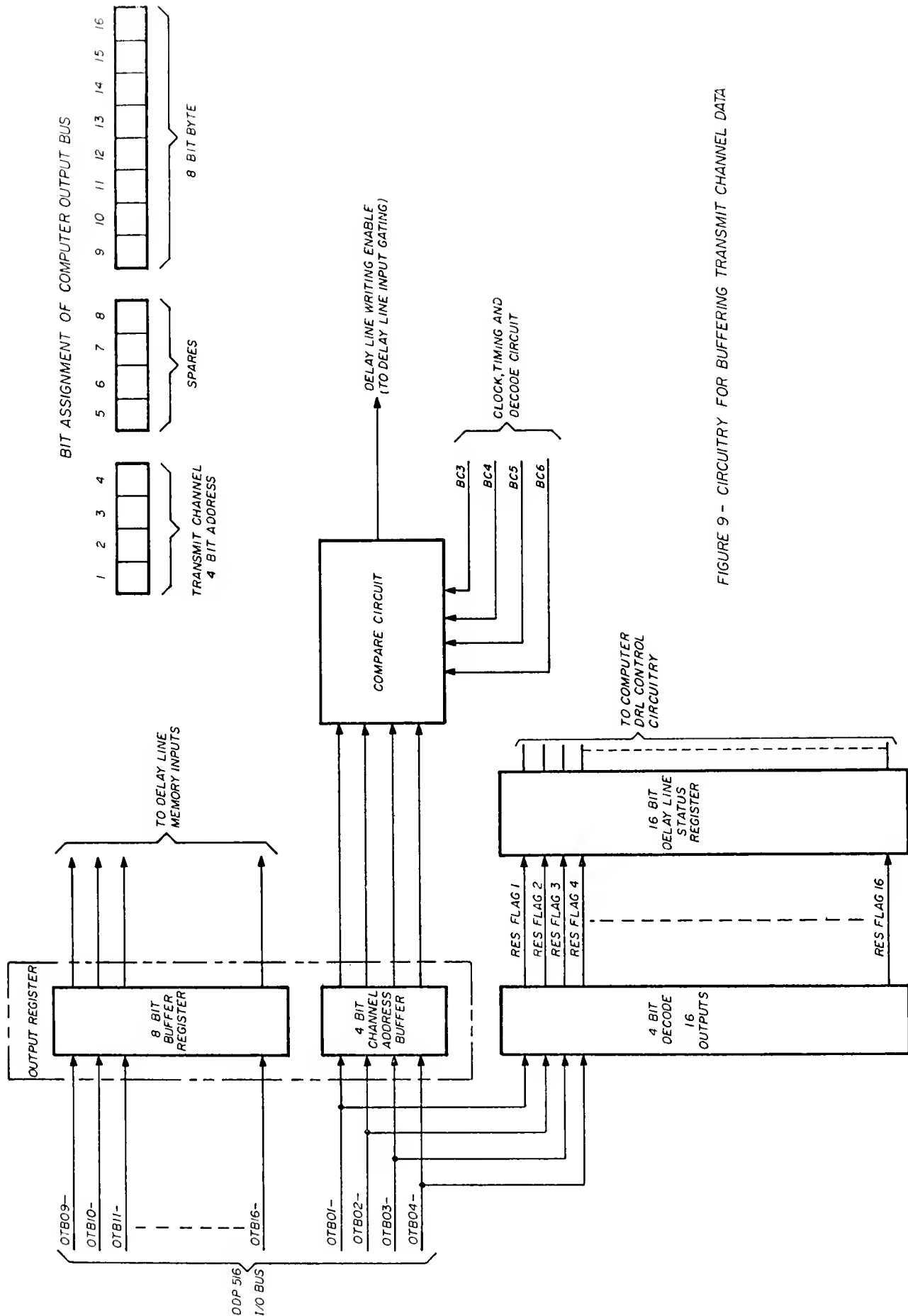


FIGURE 9 - CIRCUITRY FOR BUFFERING TRANSMIT CHANNEL DATA

## 4. SOFTWARE

### 4A. Software Specifications--Honeywell 16-Bit Series Machine

#### Interrupt Operation

SKS/430. Skip if not interrupting for input.

SKS/431. Skip if not interrupting for output.

SMK/020. Set mask for interrupting. A "1" in bit 3 of the A-register allows the multiplexer to interrupt for inputs. A "1" in bit 4 allows the multiplexer to interrupt for outputs. The master clear button on the computer console resets these bits.

#### Input Data to Computer

INA/1030. The program will skip the next instruction if the multiplexer is ready with data at the computer I/O bus. The next instruction is executed if the multiplexer is not ready to transfer data to the computer. The input data is in the format shown in Table III. Byte synchronization is performed in the multiplexer hardware. This feature can be bypassed on any individual channel if desired.

SKS/330. Skip if there are more inputs. If there are more inputs, the program may loop up to a maximum of 64  $\mu$ seconds on the next INA/1030 instruction.

#### Outputing Data from the Computer to the Multiplexer

OTA/031. The program will skip the next instruction if the multiplexer accepts the output data. If it cannot accept the data,

the next instruction in the program is executed. The output data is in the format shown in Table IV. When the OTA/031 instruction is executed, the channel delay line byte buffer and the interface register of the multiplexer must be empty in order for the output data to be accepted.

SKS/331. Skip if there are any channel byte buffers empty.

SKS/231. Skip if the interface register is empty. The program may loop to a maximum of 64  $\mu$ seconds, since the multiplexer hardware will empty the register within this time.

SKS/531. Specifying bits 1 through 4 of the output data format (Table IV) provides the test for the empty/full status of the channel delay line byte buffer. This instruction is a skip if empty.

#### Instructions for Setting Multiplexer Modes

OTA/131. This command unconditionally outputs a 16-bit word to the multiplexer. Bits 1 through 16 of the A-register represent channels 1 through 16. A "1" bit switches on the channel; a "0" bit switches off the channel. The master clear button on the computer console turns on all channels.

OTA/230. This command unconditionally outputs a 16-bit word to the multiplexer. Bits 1 through 16 of the A-register represent channels 1 through 16. A "1" bit sets the channel in the transmit mode; a "0" sets the channel in the receive mode. The master clear button on the computer console sets all channels to the receive mode.

INA/1130. This command unconditionally inputs a 16-bit word from the multiplexer. Bits 1 through 16 represent channels 1 through 16. A "1" bit indicates that a valid user is connected to the channel. A "0" bit indicates that there is no valid user connected to the channel.

Table III -- Input Data Format

<u>Bits</u>	<u>Function</u>
1 through 4	4-bit address.
5	Overwrite bit. A "1" bit indicates that the computer is not satisfying the channel input real time requirements.
6	Underwrite bit. A "1" bit indicates that the computer is not satisfying the channel output real time requirements.
7	Start bit. A "1" bit indicates the first byte of the receive message.
8	Stop bit. A "1" bit indicates that the channel has stop receiving.
9 through 16	8 bit receive byte.

Table IV -- Output Data Format

<u>Bits</u>	<u>Function</u>
1 through 4	4-bit channel address
5 through 8	Spare bits
9 through 16	8-bit transmit byte

## 4B. Computer I/O Examples

### Multiplexer Input

The delay line memory is a sequential buffer which takes 64  $\mu$ seconds to step through a complete cycle of memory. There is only one register between the delay line memory and the computer. If data can be accepted fast enough, the computer can do an effective scan of the delay line memory in 64  $\mu$ seconds.

Straight line coding is an approach that can be used for an input scheme. In order that the computer completely scan the delay line memory for inputs, the three step input instruction set (DDP-516 computer) is repeated 22 times.

INPUT	BSZ 1
	INA /1030
	JMP *+2
	STA BUFF 1
	-----
	INA /1030
	JMP *+2
	STA BUFF 2
	-----

```

      INA  /1030
      JMP  *+2
      STA  BUFF 3

```

-----  
 :

```

      INA  /1030
      JMP  *+2
      STA  BUFF 22
      JMP  * INPT

```

### Multiplexer Output

When executing an OTA/031 instruction, the first 4 bits of the computer A-register specify which of the 16 channels the output data is for. In order for this instruction to be successfully executed, the output register (computer/multiplexer interface) and the byte static buffer (delay line memory) must be empty.

Again, as for inputing data to the computer, there is an access problem. With the first 4 bits in the A-register, the SKS/531 instruction can be used to check the status of the delay line memory for a given transmit channel. Since the computer will take 6.8  $\mu$ seconds (DDP-516 computer) for each transmit channel output, data should not be outputted for two adjacent channels consecutively. This is because a word can be written into the delay line memory every 4  $\mu$ seconds.



A program that can be used for outputting data to the multiplexer follows. It is continued until all channels that require output data have been tried.

```

OUT          BSZ  1
             LDA  Chan 1
             SKS  /531
             JMP  *+4
             OTA  /031
             JMP  *-1
             JMP  *+2
             STA  SAV 1
-----

             LDA  Chan 3
             SKS  /531
             JMP  *+4
             ØTA /031
             JMP  *-1
             JMP  *+2
             STA  SAV 3
-----

Channel 5 data routine
-----
             :
             :
             :
-----

Channel 15 data routine
-----

Channel 2 data routine
-----
             :
             :

```

⋮

---

Channel 16 data routine

---

## 5. APPENDICES

### Appendix A. RS-232-B Serial Channel Standard

The RS-232-B standard serial data channel is terminated by a 25 pin Cinch connector. Signals on the input and output pins are bi-polar DC levels: -6 volts or +6 volts. The connector pins and their corresponding channel functions are as follows:

Pin 1--Protective Ground (AA)

Pin 2--Transmitted Data (BA). The terminal supplies the NRZ data stream to this pin. A logical "1" is -6 volts. A logical "0" is +6 volts.

Pin 3--Received Data (BB). This pin outputs the NRZ data stream to the terminal. A logical "1" is -6 volts. A logical "0" is +6 volts.

Pin 4--Request to Send (CA). When transmitting data onto the channel, the terminal equipment must supply +6 volts to this pin. When receiving data from the channel, the terminal equipment must supply -6 volts to this pin.

Pin 5--Clear to send (CB). This pin supplies a +6 volt transmit enable signal to the terminal. A -6 volt signal will inhibit the terminal from transmitting. When a

terminal issues a CA signal to the channel, the channel will return, after a delay, a CB signal to the terminal.

Pin 6--Data Set Ready (CC). This pin supplies a +6 volt signal to the terminal to indicate when the channel can transmit or receive data. Ground indicates that the channel is not ready for communicating.

Pin 7--Signal Ground (AB).

Pin 8--Data Carrier Detector (CF). This pin supplies +6 volts to the terminal to indicate the presence of the Receive Clock (DD). -6 volts indicates the absence of DD.

Pin 9--Reserved for data set testing.

Pin 10--Reserved for data set testing.

Pin 11--Unassigned.

Pin 12--Unassigned.

Pin 13--Unassigned.

Pin 14--Unassigned.

Pin 15--Channel Transmit Clock (DB). This pin supplies a +6 volt square wave to the terminal equipment when transmitting data into the channel. The positive transition of DB is used for shifting data on to the BA line from the terminal.

Pin 16--Unassigned.

Pin 17--Receiver Clock (DD). This pin supplies a +6 volt square wave to the terminal equipment when receiving data from the channel. The BB line presents a new data bit at the

positive transition of DD.

Pin 18-Unassigned.

Pin 19-Unassigned.

Pin 20-Data Terminal Ready (CD). The terminal supplies +6 volts to this pin to indicate it is ready to operate. -6 volts indicates a not ready condition.

Pin 21-Unassigned.

Pin 22-Ring Indicator (CE). This pin supplies the terminal with +6 volts to indicate that there is a remote user trying to establish a communications link through the channel. Otherwise this pin is at -6 volts.

Pin 23-Unassigned.

Pin 24-External Transmit Clock (DA). The terminal supplies to this pin a  $\pm 6$  volt square wave when the transmitted data through the channel is to be clocked by the terminal equipment.

Pin 25-Unassigned.

## Appendix B. Honeywell 16-Bit Series Standard I/O Channel

The Honeywell 16-bit series standard I/O channel is a shared bus I/O system that consists basically of 48 twisted lead pairs. Thirty pairs are used for output and 18 pairs for input.

The 30 output twisted lead pairs are:

- 16-for data bits (OTBXX)
- 10-for address bits (ADBXX)
- 1-for reset ready (RRL)
- 1-for output command (OCP)
- 1-for master clear (MSTCL)
- 1-for set mask (SMSK)

The 18 input twisted lead pairs are:

- 16-for data bits (INBXX)
- 1-for device ready (DRL)
- 1-for program interrupt (PIL)

Signals on the output and input leads are DC levels: 0 or -6 volts in the DDP-116 and 0 or +6 volts in the DDP-416, DDP-516, and H-316. All leads are shared with all interfaces. Memory access I/O channels are available as options for all these 16-bit computers

The channel is synchronized by a specific time sequence of events. For example, within the I/O instructions there exist the following sequence of events: (a) an address goes out on the address bus, (b) a device decodes this address and responds, and (c) a strobe for the device terminates the instruction. The standard channel provides a hardware program interrupt feature to service devices on demand. The standard channel is a parallel I/O bus that is shared by all interfaces under program control.

The I/O software instructions and their hardware significance are discussed briefly.

#### Output Command (OCP)

This instruction initiates DC levels on the 10-bit address bus and a delay pulse on the output command twisted pair. This instruction is used to control the functioning of the I/O device, and expects no response from the device.

#### Sense Condition (SKS)

This instruction initiates DC levels on the 10-bit address bus and senses the input device ready lead. When a signal appears on this lead, during a set time in the instruction period, it causes the next instruction to be skipped. If no signal appears, the next instruction will be executed. This (SKS) instruction is used to test conditions of specific interfaces on the I/O bus.

### Single Word Input (INA)

The input instruction initiates action by placing DC levels on the 10-bit address bus. The addressed interface then places its input data on the 16-bit input bus; simultaneously, a signal is placed on the input device ready lead. The computer will acknowledge the input with a signal on the output reset ready lead, and will skip the next instruction to be executed. In the event of no input, the addressed interface does not return a device ready signal. The computer does not skip the next instruction, nor acknowledge an input with a signal on the reset ready lead. Multiple word inputs must be accomplished by creating a program loop of single word (INA) inputs.

### Single Word Output (ØTA)

The output instruction initiates action by placing DC levels on the 10-bit address bus and the 16-bit output bus. The addressed interface accepts the data by placing a signal on the input device ready lead. The computer will then respond on the reset ready lead with a signal which can be used as a strobe for the output data. The computer will then skip the next instruction to be executed. In the event the addressed interface does not accept the output, the computer will not skip the next instruction, nor respond with a signal on the reset ready lead. Multiple word outputs must be accomplished by creating a program loop of single word (ØTA) outputs.



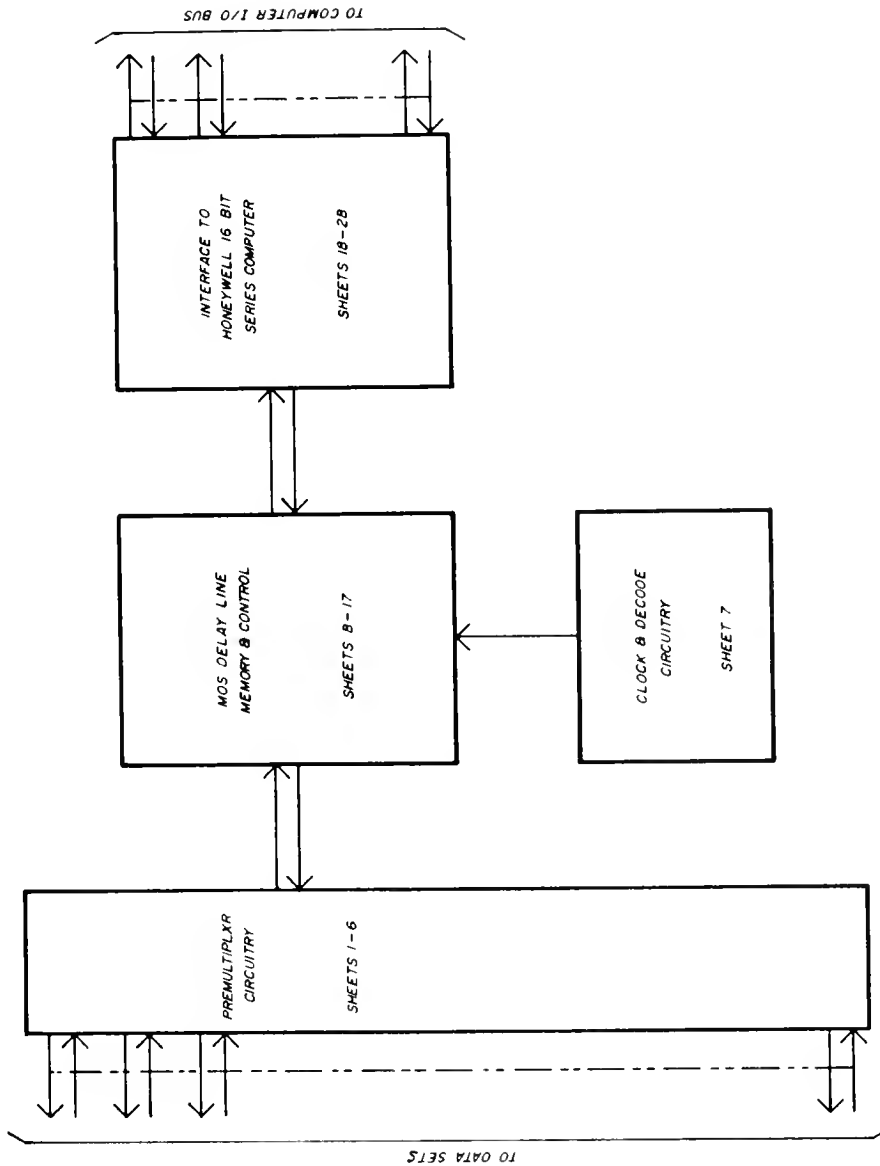
Any interface on the I/O bus can initiate an interrupt by placing a DC level on the program interrupt lead (PIL). Since this lead is common to all interfaces, when an interrupt occurs a program using the sense condition instruction (SKS) must be used to seek out the source of interrupt, and then jump to a subroutine to service it. Priorities of program interrupts are accomplished through the use of mask bits. Each interface is assigned a bit which is used to generate its program interrupt (PIL) signal.

Refer to Honeywell Interface Manual (Pub. No. 130071624) for further details on the I/O channel of the 16-bit series computers.



## Appendix C. System Schematics





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28	DRL CIRCUIT
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TITLE  
16 PORT DATA SET MULTIPLXR



TABLE ONE

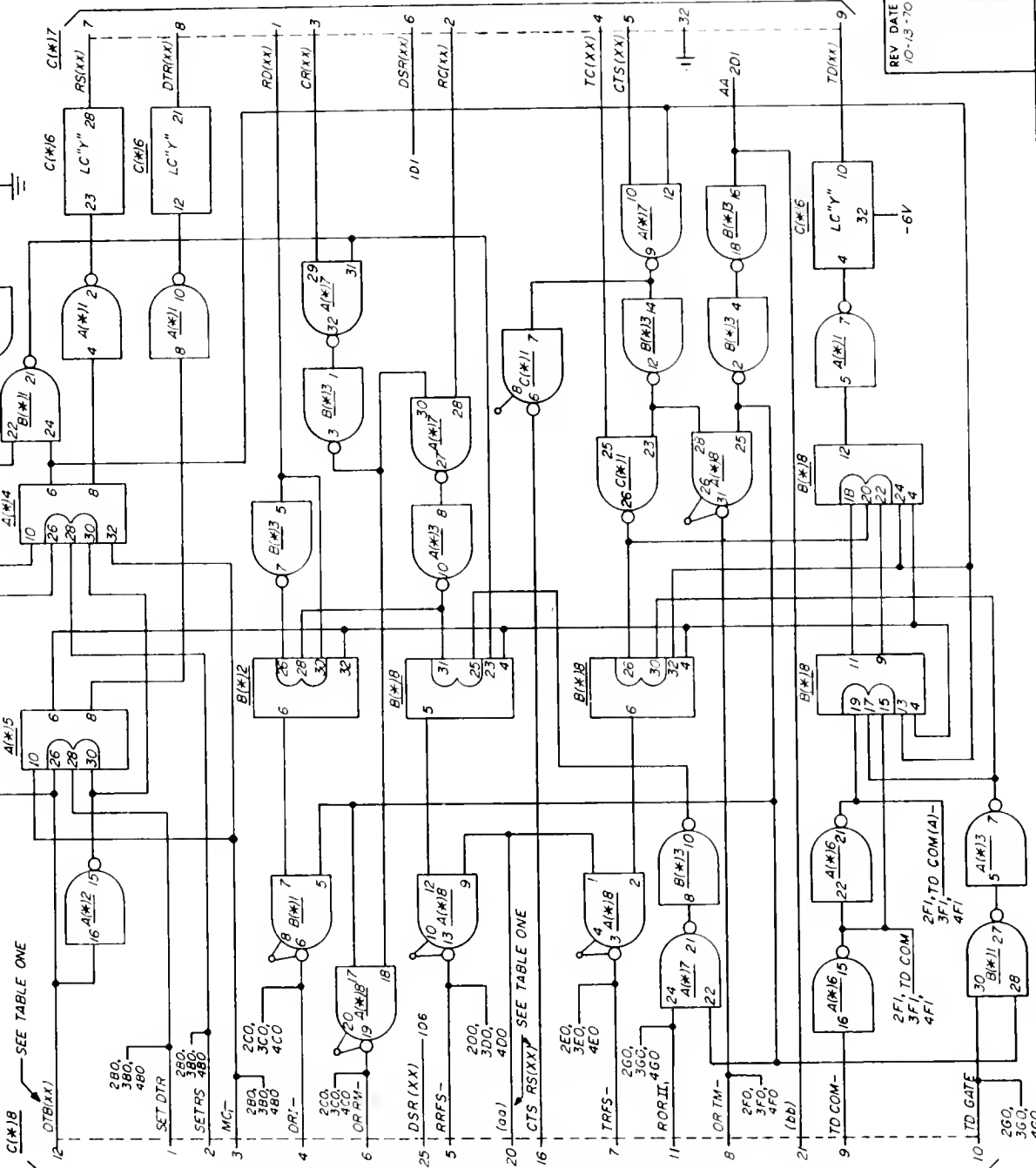
CKT	(ac)	(ad)	(bb)	(xx)	(*)	(ab)
1	A28	17	00	01	2	A17
5	A26	03	04	05	3	C12
9	A24	07	10	09	4	A67
13	A22	13	14	13	5	C62

TO LEVEL CONVERTOR BLOCK  
CONNECT LL-PAC CABLE 2, (SHT 5)  
TO POSITION (ab), SEE TABLE ONE

NOTE ALL EQUIPMENT ON THIS SHEET  
IS ON RACK NO 2

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TITLE  
PREMULTIPLYR 1,5,8/13



TO HUB RACK NO 1  
(CONNECT LL-PAC CABLE TO POSITION (ac), SEE TABLE ONE) A SHT 6, FO









TABLE THREE

CRT	(oc)	(ob)	(db)	(xx)	(M)
3	A28	C17	O2	O3	2
7	A26	A12	O6	O7	3
11	A24	C67	I2	I1	4
15	A22	A62	I5	I5	5

TO LEVEL CONVERTOR BLOCK  
CONNECT U-PAC CABLE 'A' (SHT 5)  
TO POSITION (ob), SEE TABLE THREE

NOTE: ALL EQUIPMENT ON THIS SHEET  
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TITLE  
PREMULTIPLXR 3.7/11B/5





CAK	(oc)	(ob)	(ob)	(Xx)	(*)
4	A28	C17	03	04	2
8	A26	A12	07	08	3
12	A24	C67	13	12	4
16	A22	A62	17	16	5

TO LEVEL CONVERTOR BLOCK  
CONNECT U-PAC CABLE 'A' (SHT 5)  
TO POSITION (46), SEE TABLE FOUR

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DRAWING NO 540		SHEET NO 4			

TITLE
PREMULTIPLY XR 4 8 12 8 16



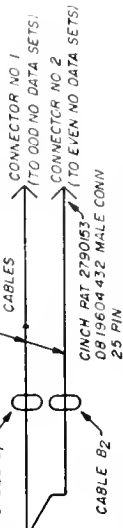
TABLE FIVE

DATA SET	(xx)	(xy)	(aq)	(bb)	(cc)	(dd)	(ee)	(ff)	(gg)	(hh)	(ij)
1	01	1	C27 A17	A16	A15	—	—	B18	B17	A18	2
2	02	2	C27 A17	—	—	C15	—	—	C18	—	3
3	03	3	C22 C17	—	—	—	B16	B15	C18	4	—
4	04	4	C22 C17	—	—	C14	—	—	C11	—	5
5	05	1	C37 C12	—	—	—	B13	B14	C11	6	—
6	06	2	C37 C12	—	—	A14	—	—	A11	—	7
7	07	3	C32 A12	A13	A14	—	—	B11	B12	A11	8
8	08	4	C32 A12	—	—	—	—	—	—	—	9
9	09	1	C47 A67	A66	A65	—	—	B66	B67	A68	10
10	10	2	C47 A67	—	—	C65	—	—	C68	—	11
11	11	3	C42 C67	—	—	—	B66	B65	C68	12	—
12	12	4	C42 C67	—	—	C64	—	—	C61	—	13
13	13	1	C57 C62	C63	C64	—	—	B63	B64	C61	14
14	14	2	C57 C62	—	—	—	—	—	—	—	15
15	15	3	C52 A62	A63	A64	—	—	B61	B62	A61	16
16	16	4	C52 A62	—	—	—	—	—	—	—	—

TO ODD NUMBER  
DATA SET (ij)TO EVEN NUMBER  
DATA SET (hh)

TO PRE-MULTIPLEXOR

TO DATA SETS



TERM NO. LL-PAC	SIGNAL	DATA SET CONN
1	RD--	3
2	RC--	17
3	CR--	8
4	TC--	15
5	CTS--	5
6	DSR--	6
7	RS--	4
8	DTR--	20
9	TO--	2
32	GRD	7 (STRAP 7B1)
11	RD--	3
12	RC--	17
13	CR--	8
14	TC--	15
15	CTS--	5
16	DSR--	6
17	RS--	4
18	DTR--	20
19	TO--	2
33	GRD	7 (STRAP 7B1)

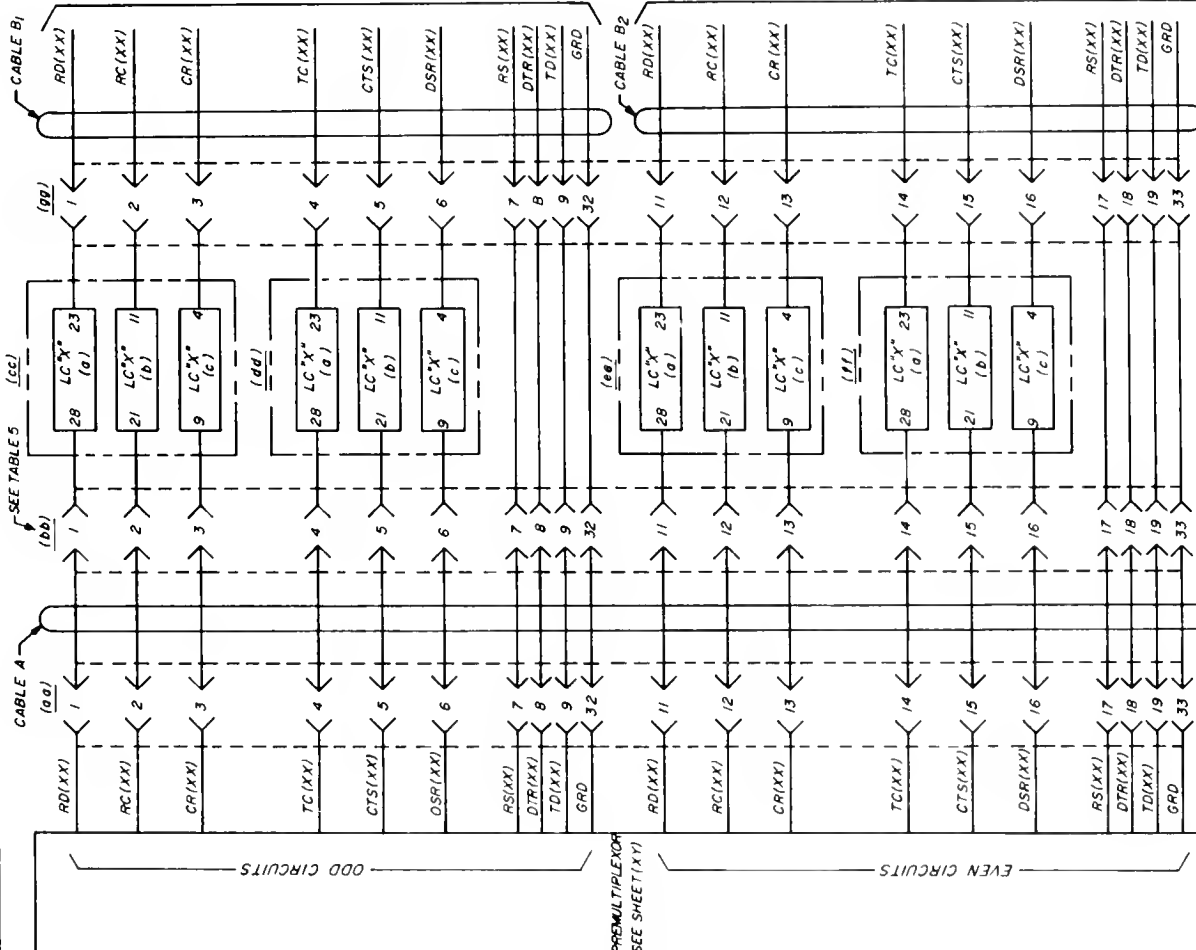
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 DATE 7-10-70  
 TITLE PREMULTIPLX TO DATA SET CABLING
PREMULTIPLX  
SEE SHEET (xy)





TABLE 6

'A <sub>6</sub> ' PRE-MULT LOCATION	'B <sub>6</sub> ' SIGNAL	'C <sub>6</sub> ' PIN	'D <sub>6</sub> ' HUB LOCATION
		35	
		34	
		33	
		32	
		31	
		30	
		29	
400	DSR 16	28	23F4
300	DSR 15	27	23E4
200	DSR 14	26	23E4
100	DSR 13	25	23D4
400	17	24	7C9
300	16	23	7C9
200	15	22	7C9
100	14	21	7B9
400	13	20	7B9
300	12	19	26G4
200	11	18	26E4
100	10	17	26D4
400	9	16	26B4
300	8	15	27D0
200	7	14	27D1
100	6	13	27D1
400	5	12	27D2
300	4	11	809
200	3	10	8C9
100	2	9	8F9
400	1	8	8B0
300	0	7	8F0
200		6	8A0
100		5	8E0
400		4	8G0
300		3	22D0
200		2	21C9
100		1	21B9

TABLE 7

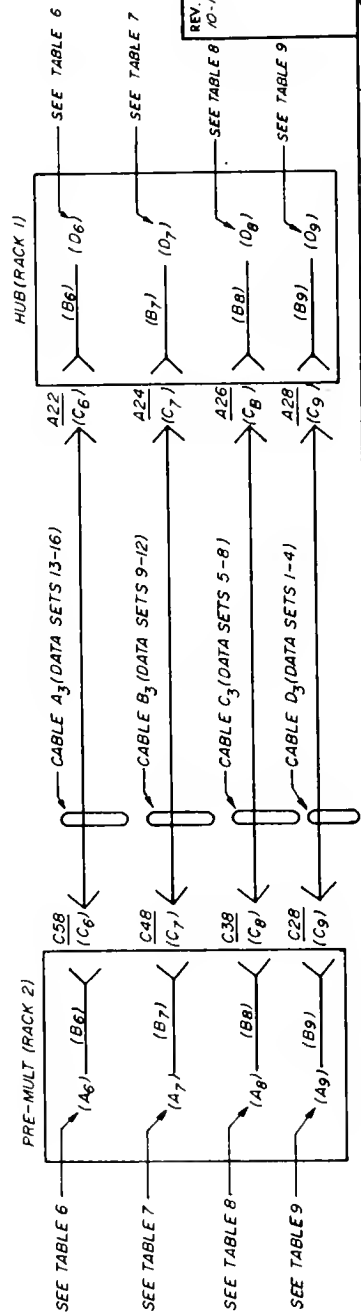
'A <sub>7</sub> ' PRE-MULT LOCATION	'B <sub>7</sub> ' SIGNAL	'C <sub>7</sub> ' PIN	'D <sub>7</sub> ' HUB LOCATION
		35	
		34	
		33	
		32	
		31	
		30	
		29	
400	DSR 12	28	23C4
300	DSR 11	27	23C4
200	DSR 10	26	23B4
100	DSR 09	25	23B4
400	13	24	7B9
300	12	23	7B9
200	11	22	7A9
100	10	21	7A9
400	07	20	7F9
300	06	19	26G0
200	05	18	26E0
100	04	17	26D0
400	03	16	26B0
300	02	15	27D3
200	01	14	27D3
100	00	13	27D4
400		12	809
300		11	8C9
200		10	8F9
100		9	8B0
400		8	8F0
300		7	8A0
200		6	8E0
100		5	8G0
400		4	22D0
300		3	21C9
200		2	21B9
100		1	21B9

TABLE 8

'A <sub>8</sub> ' PRE-MULT LOCATION	'B <sub>8</sub> ' SIGNAL	'C <sub>8</sub> ' PIN	'D <sub>8</sub> ' HUB LOCATION
		35	
		34	
		33	
		32	
		31	
		30	
		29	
400	DSR 08	28	23F1
300	DSR 07	27	23E1
200	DSR 06	26	23D1
100	DSR 05	25	23D1
400	07	24	7F9
300	06	23	7F9
200	05	22	7F9
100	04	21	7E9
400	03	20	7E9
300	02	19	25G4
200	01	18	25E4
100	00	17	25D4
400		16	25B4
300		15	27D5
200		14	27D5
100		13	27D6
400		12	27D7
300		11	809
200		10	8C9
100		9	8F9
400		8	8B0
300		7	8F0
200		6	8A0
100		5	8E0
400		4	8G0
300		3	22D0
200		2	21C9
100		1	21B9

TABLE 9

'A <sub>9</sub> ' PRE-MULT LOCATION	'B <sub>9</sub> ' SIGNAL	'C <sub>9</sub> ' PIN	'D <sub>9</sub> ' HUB LOCATION
		35	
		34	
		33	
		32	
		31	
		30	
		29	
400	DSR 04	28	23C1
300	DSR 03	27	23C1
200	DSR 02	26	23B1
100	DSR 01	25	23B1
400	03	24	7E9
300	02	23	7E9
200	01	22	7D9
100	00	21	7D9
400		20	7C9
300		19	25G0
200		18	25E0
100		17	25D0
400		16	25B0
300		15	27D7
200		14	27D8
100		13	27D8
400		12	27D9
300		11	809
200		10	8C9
100		9	8F9
400		8	8B0
300		7	8F0
200		6	8A0
100		5	8E0
400		4	8G0
300		3	22D0
200		2	21C9
100		1	21B9



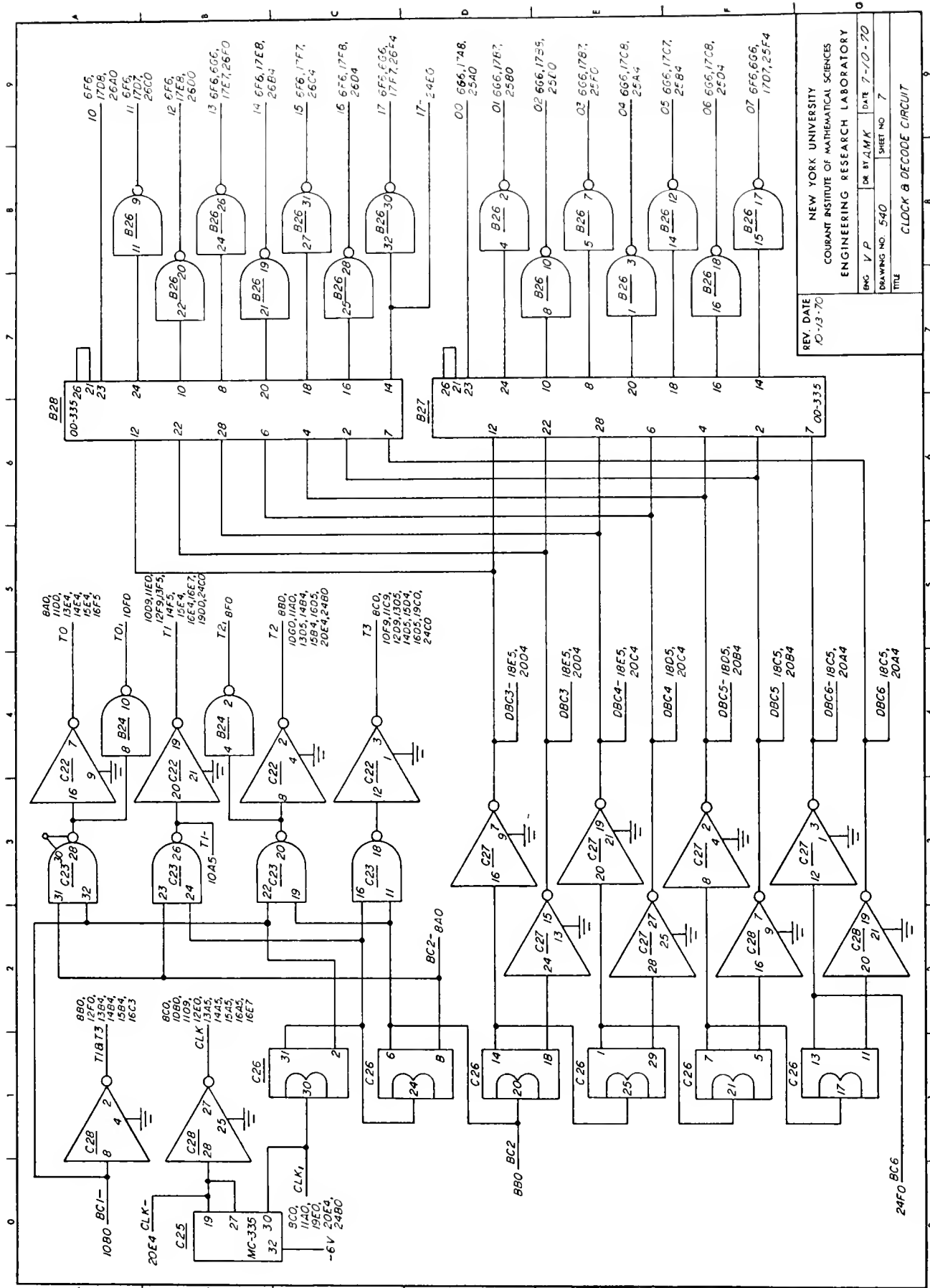
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HUB TO PREMULTIPLYR CABLING





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DWG. V P    DR. BY L.M.K.    DATE 7-10-70  
DRAWING NO. 540    SHEET NO. 7  
TITLE  
CLOCK & DECODE CIRCUIT



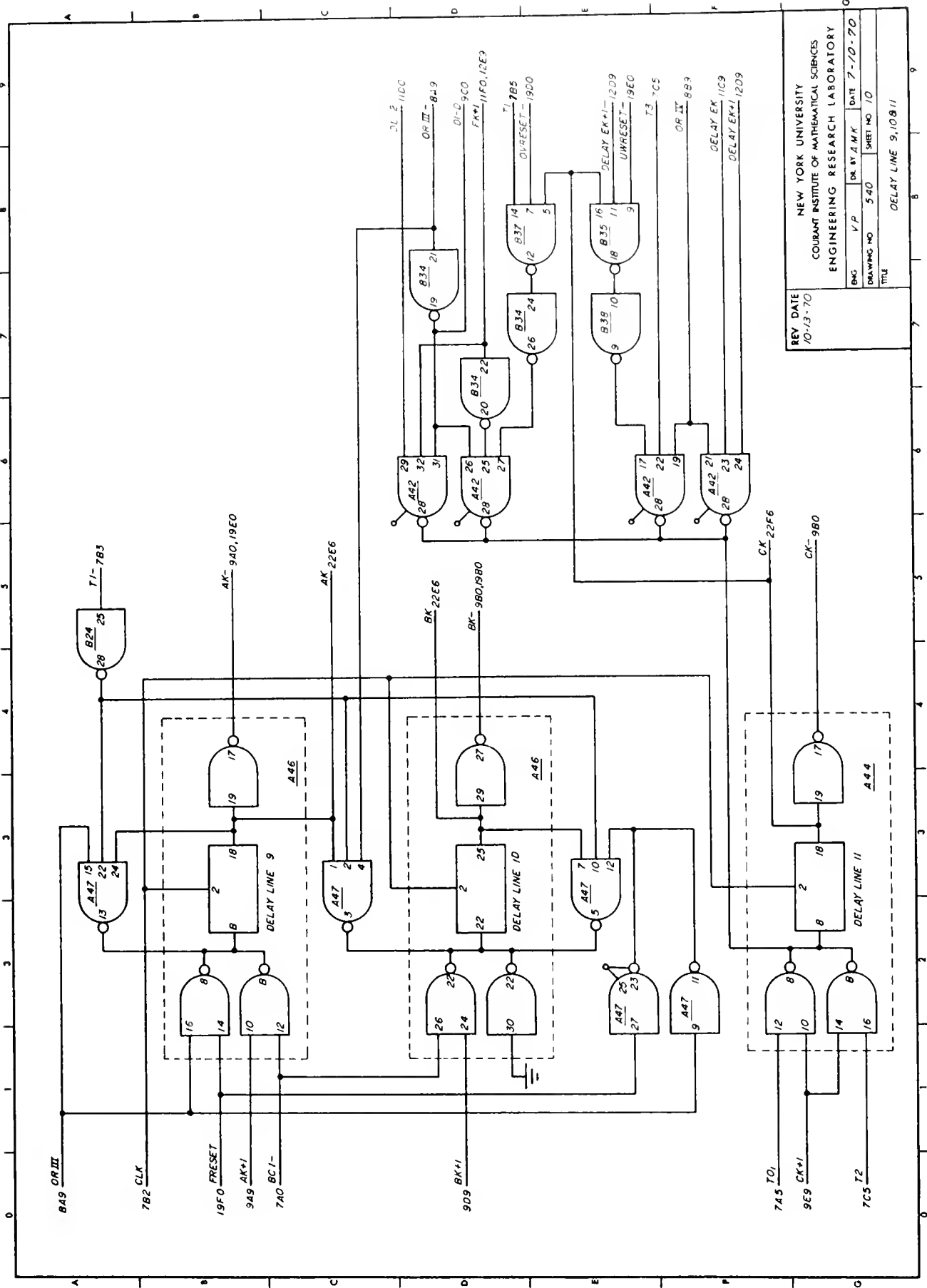










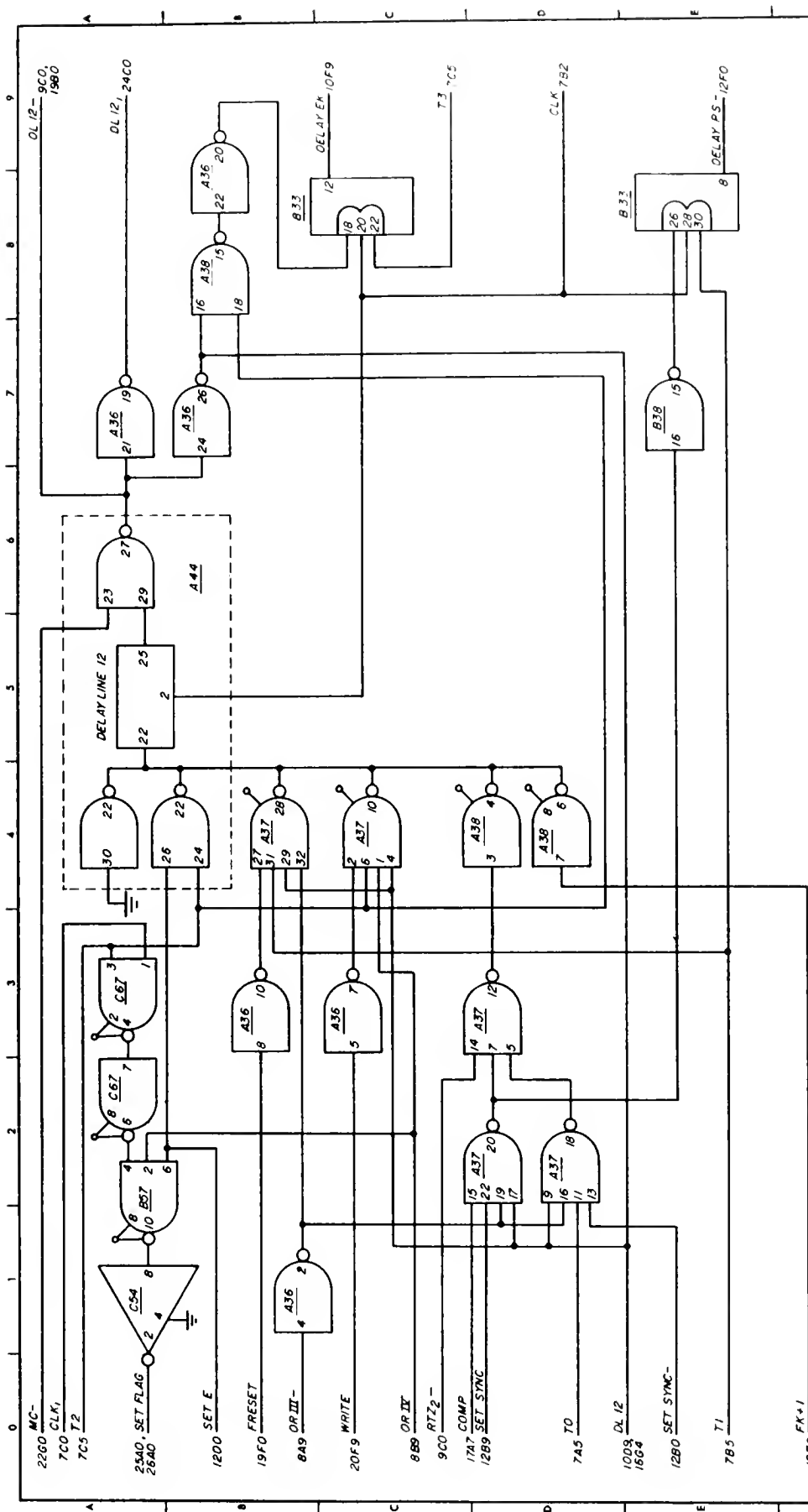


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DRAWING NO. 540 SHEET NO. 10  
TITLE  
DELAY LINE 9,10,11



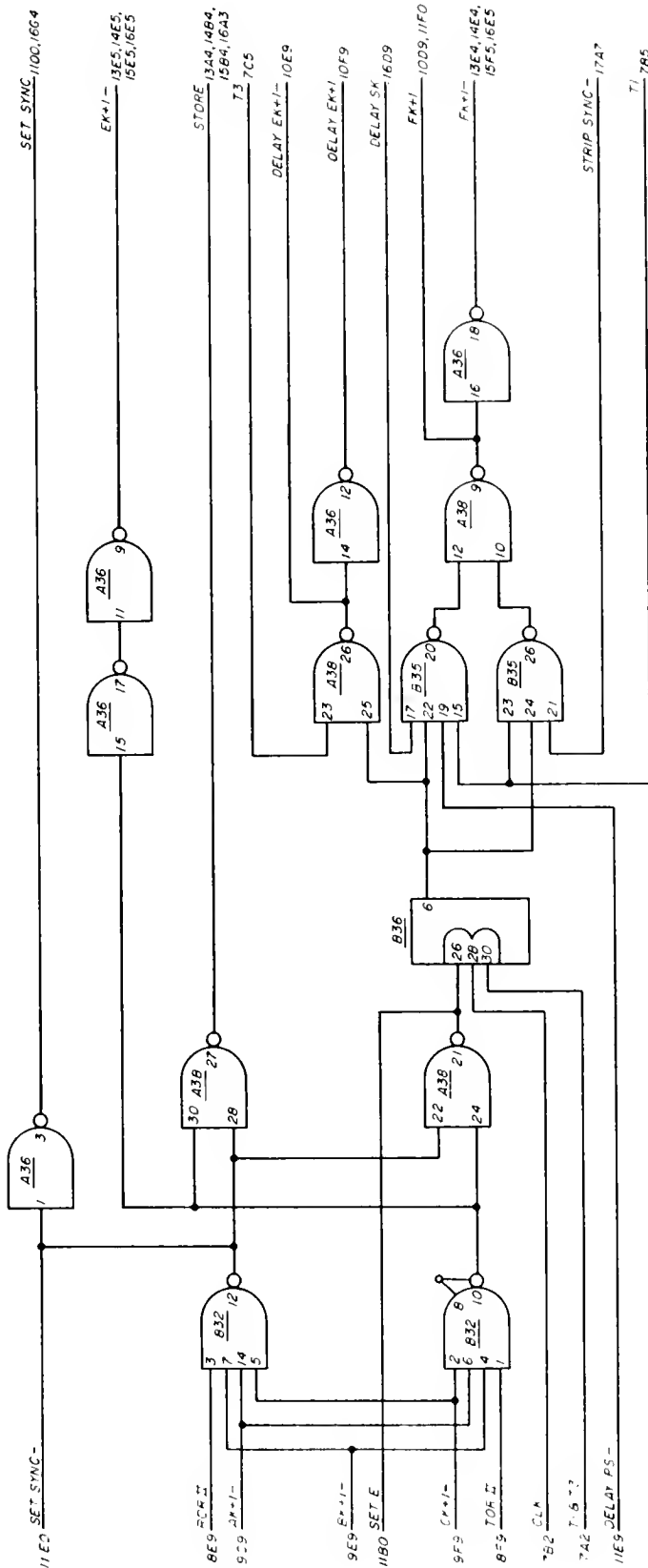


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DRAWING NO	540	SHEET NO	11	
TITLE DELAY LINE 12 B CONTROL LOGIC				





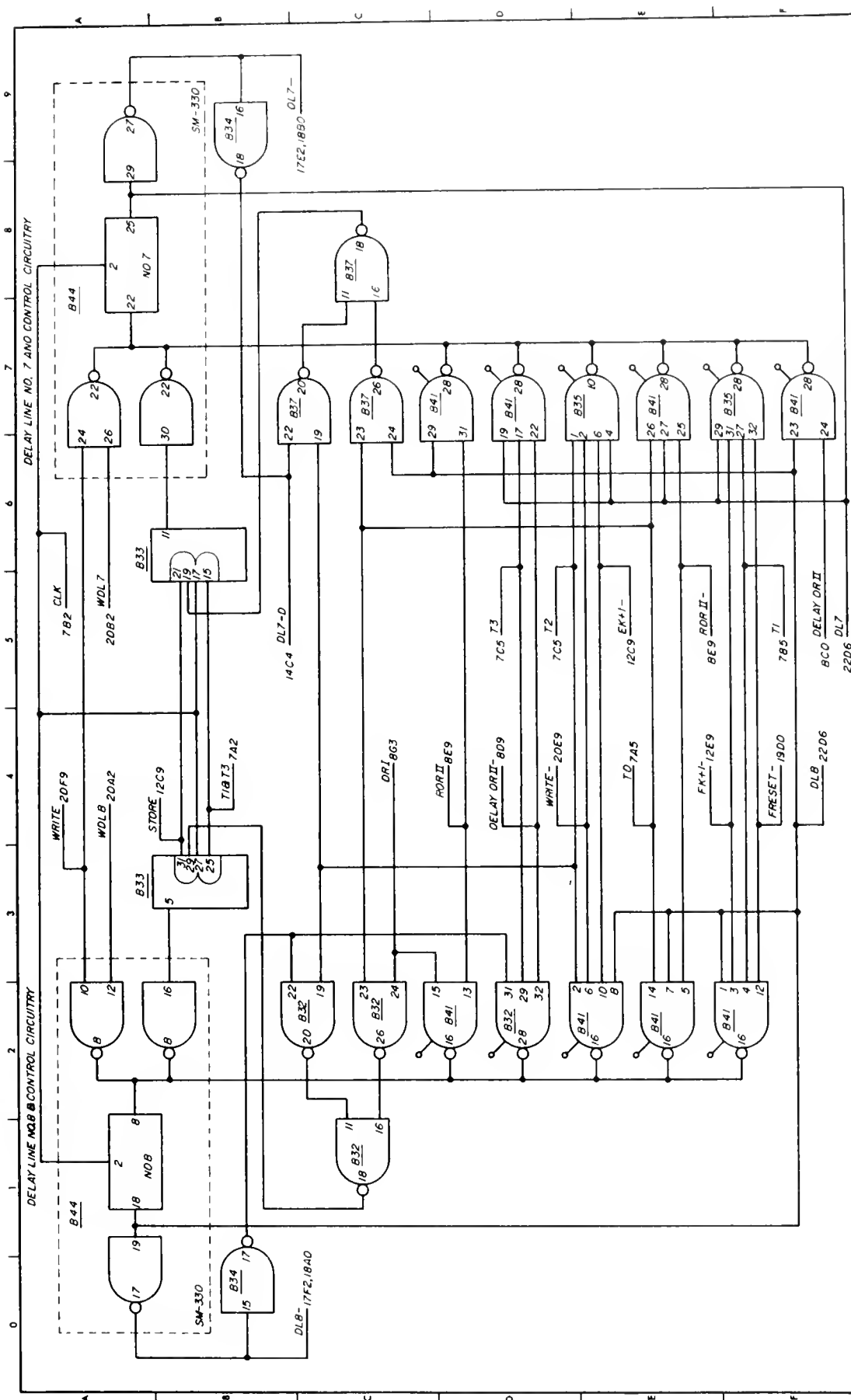
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SHEET NO. 12

TITLE  
CONTROL LOGIC FOR DELAY LINE 12





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TITLE		DELAY LINES 788	



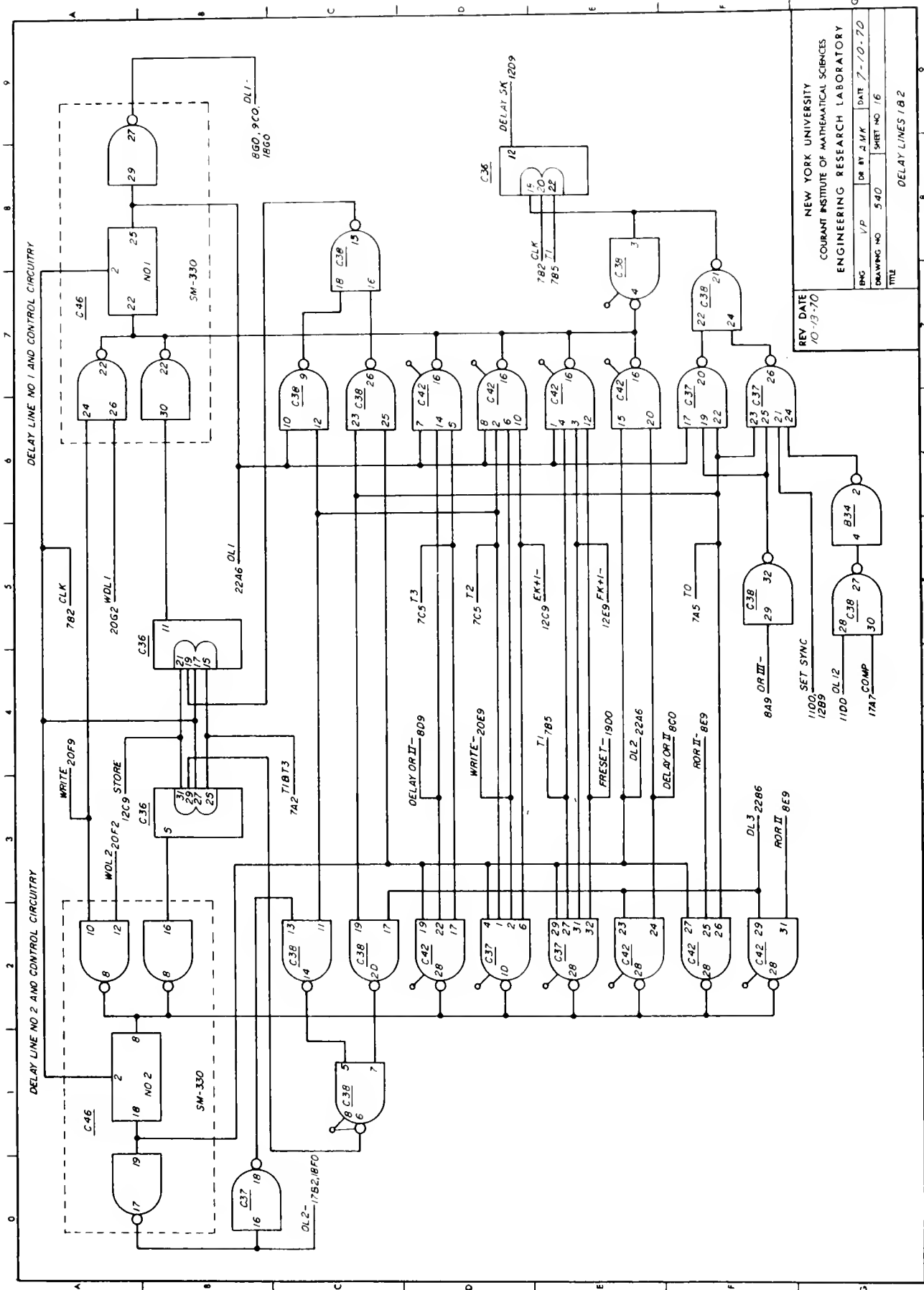












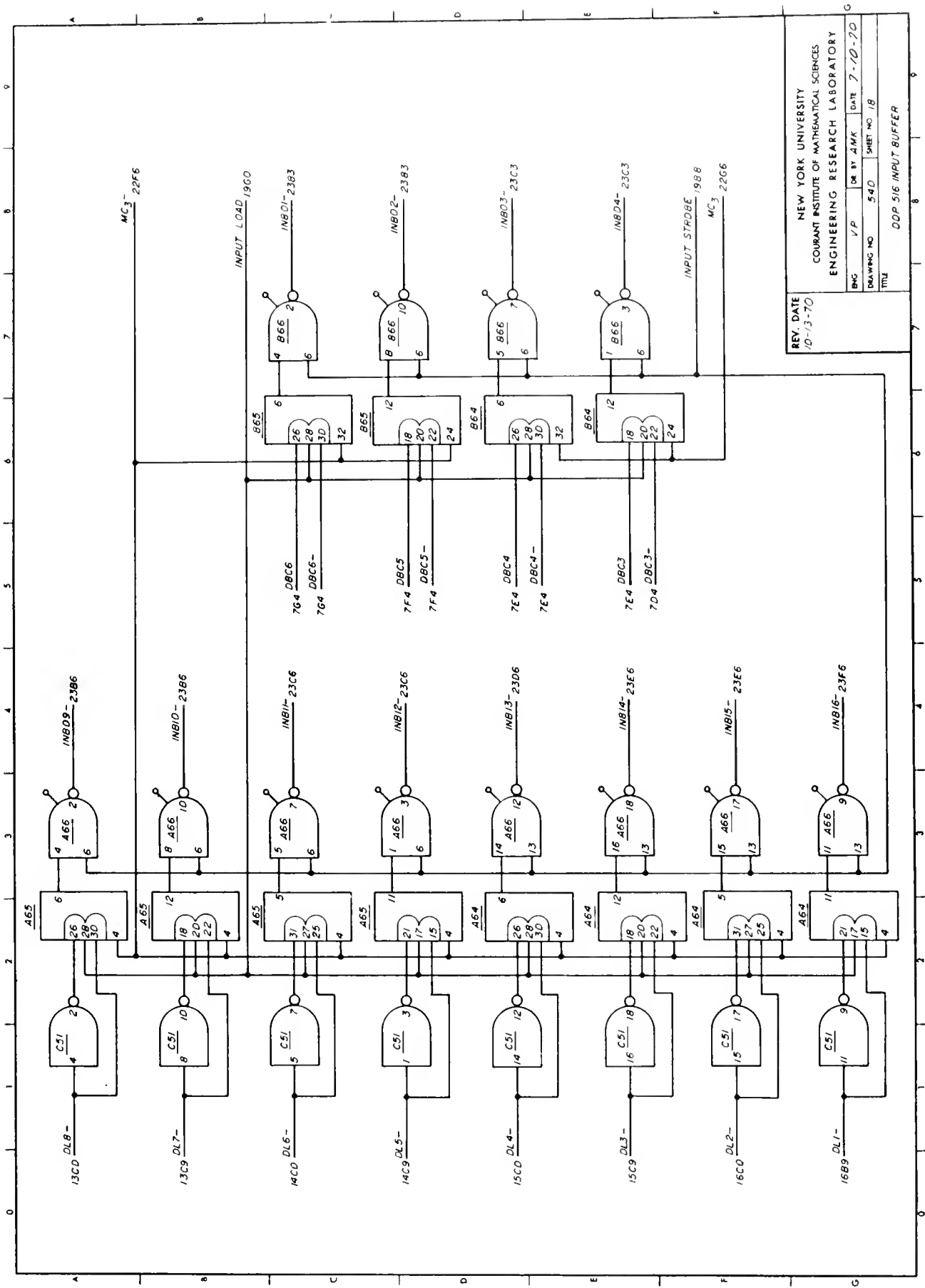
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		SHEET NO	16
		TITLE	DELAY LINES 1/8 2





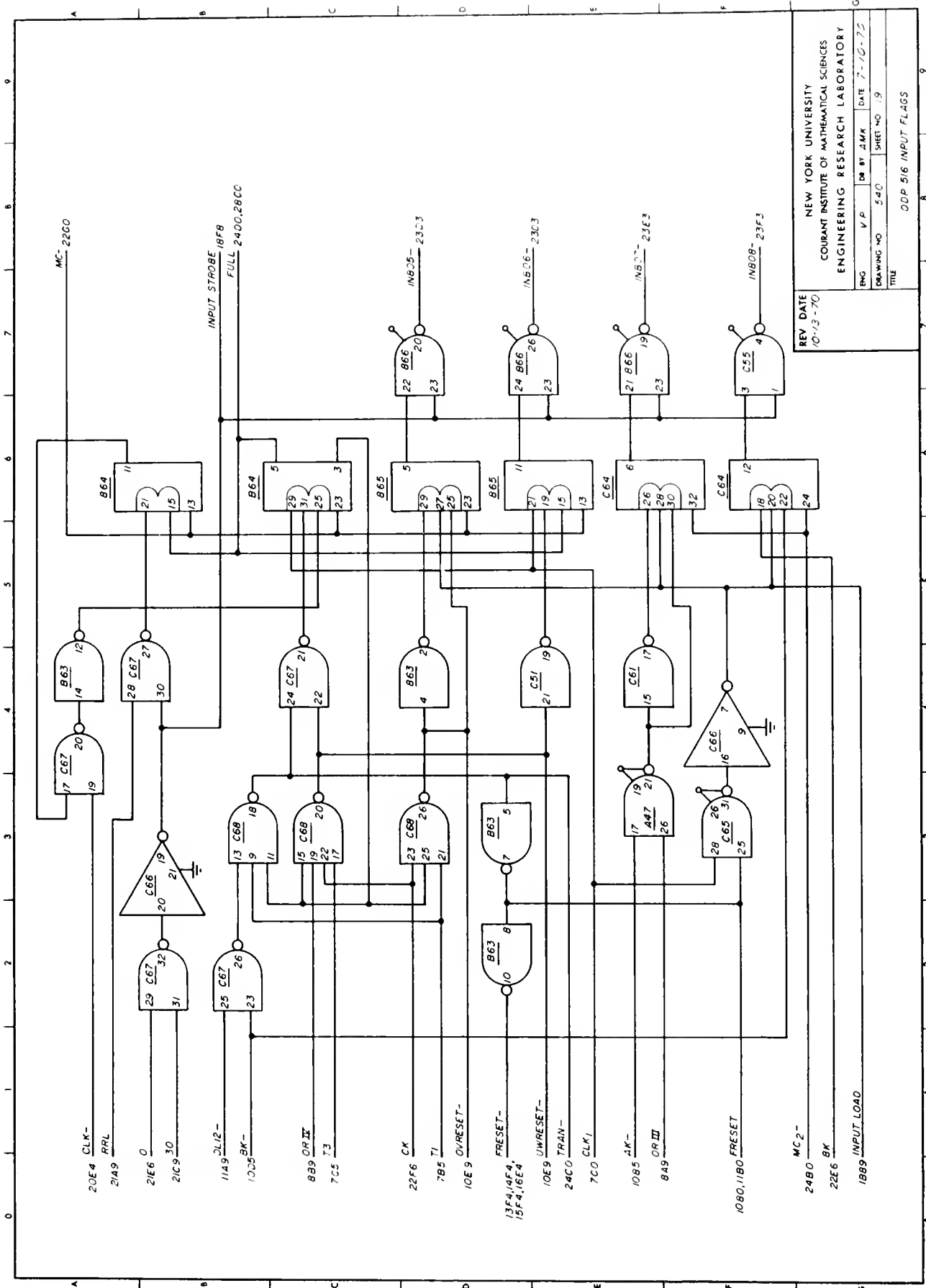






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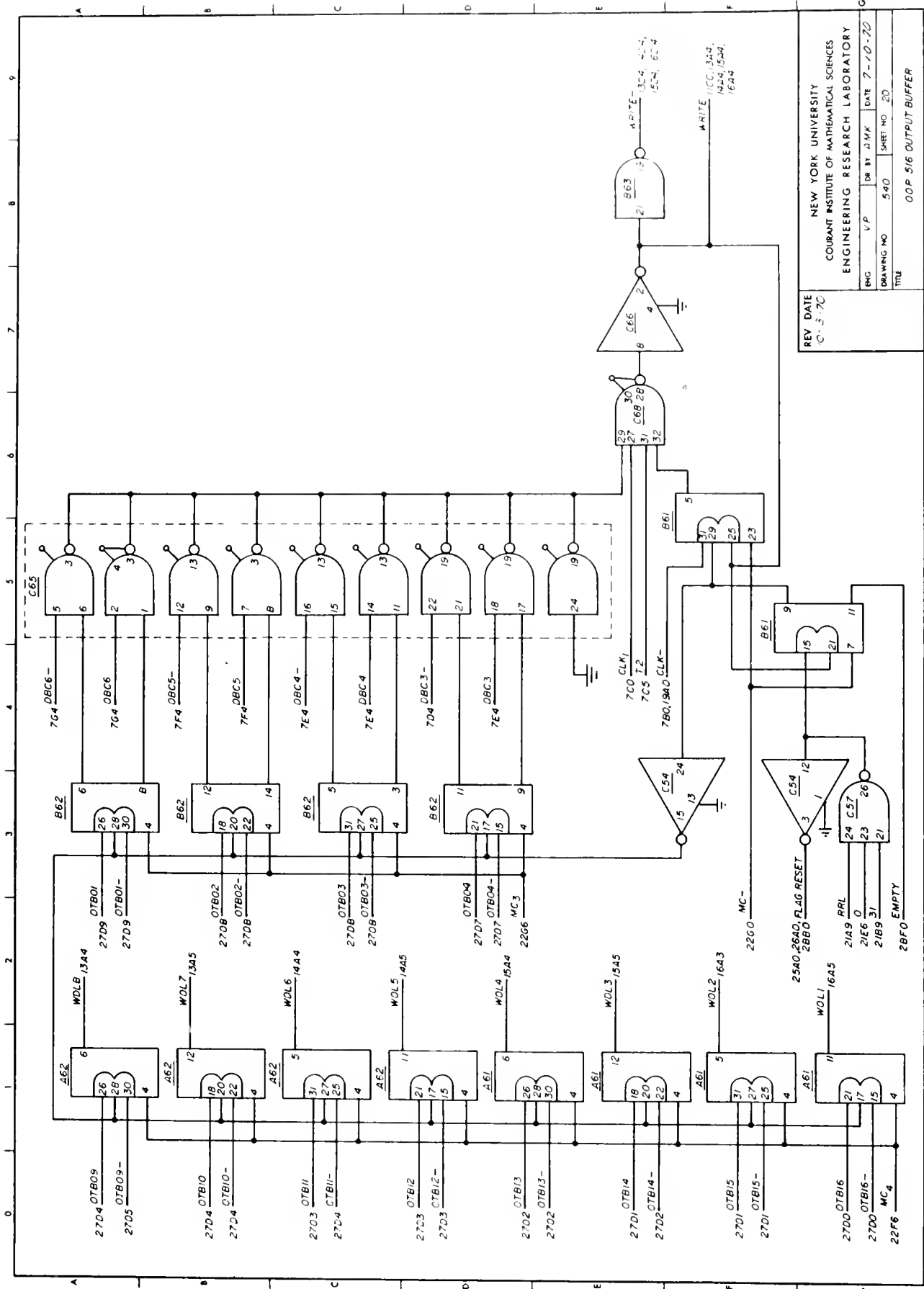


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DRAWING NO 540 SHEET NO 19  
TITLE ODP 516 INPUT FLAGS





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DRAWING NO. 540		TITLE	
		00P 516 OUTPUT BUFFER	



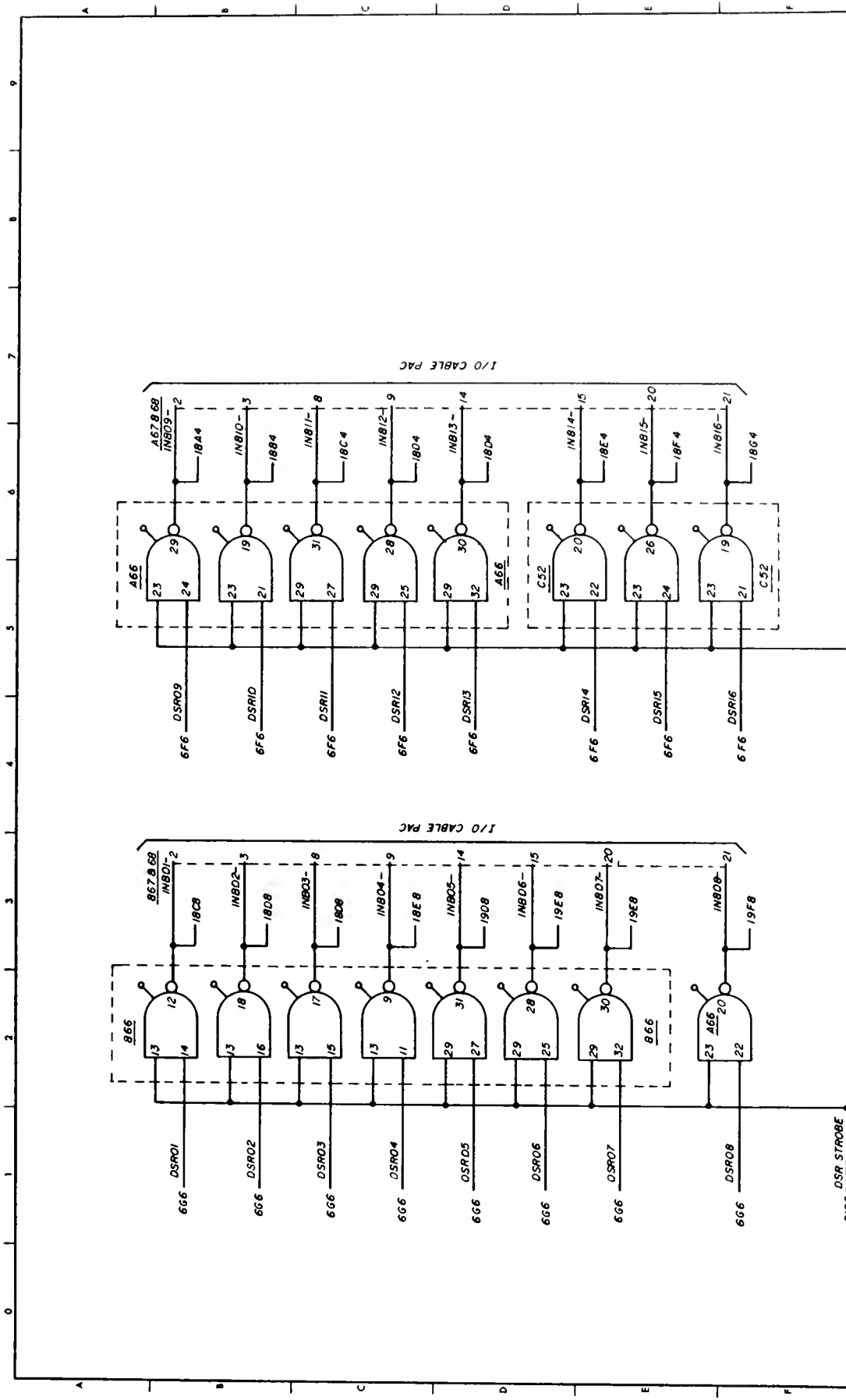








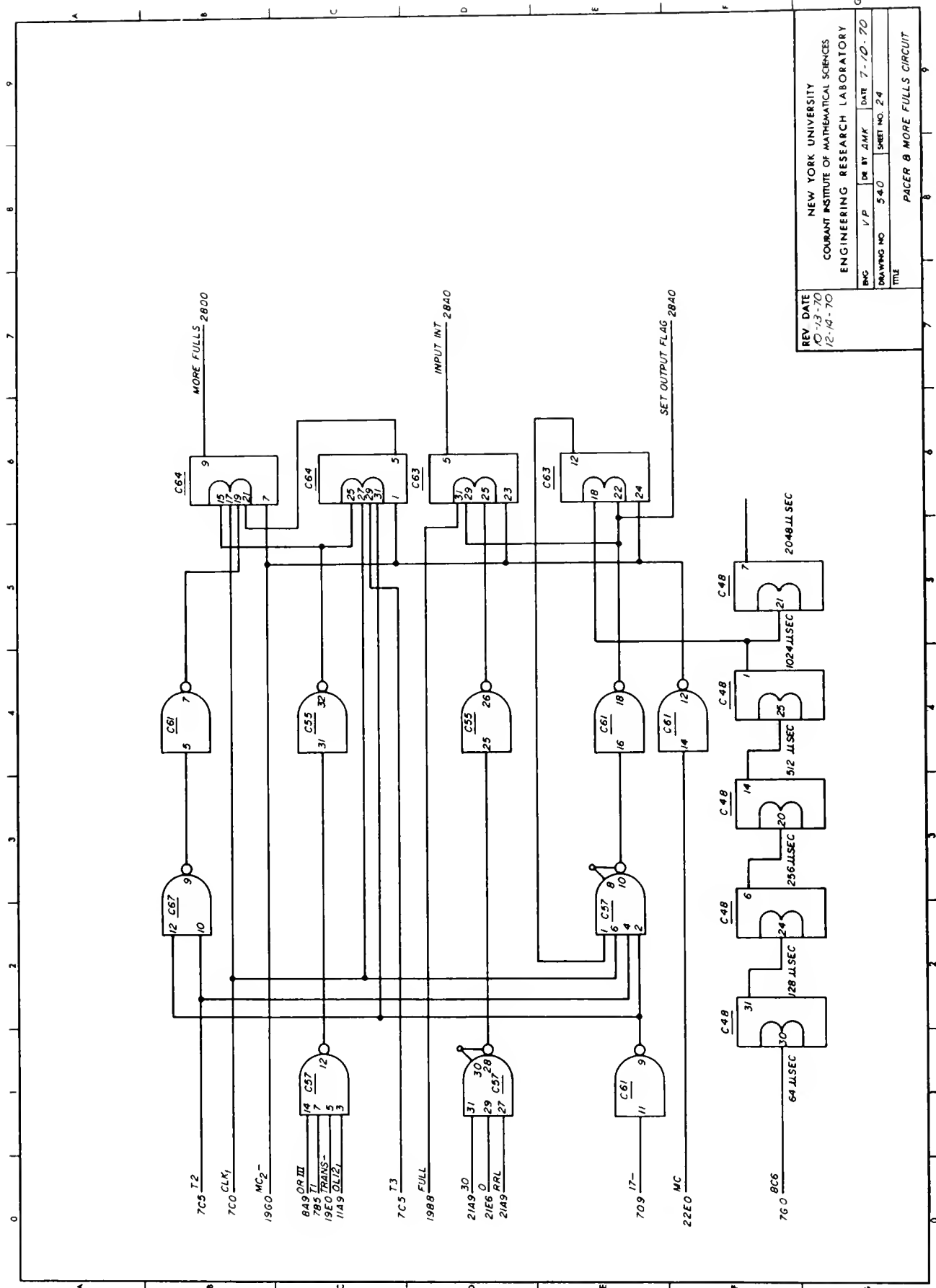




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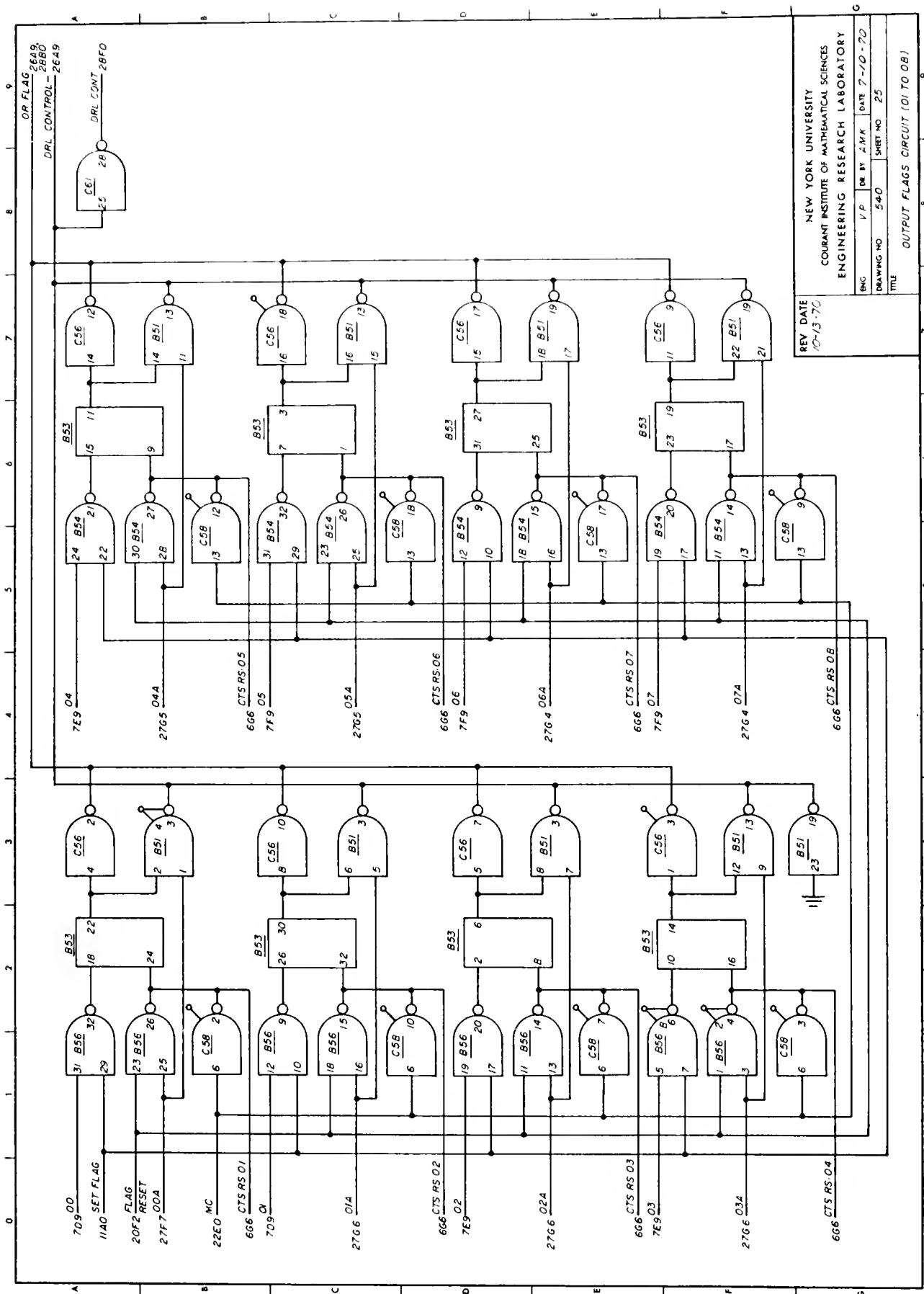
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TITLE DSR INPUT BUS			





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DRAWING NO	540		SHEET NO. 24		
TITLE	PACER 8 MORE FULLS CIRCUIT				

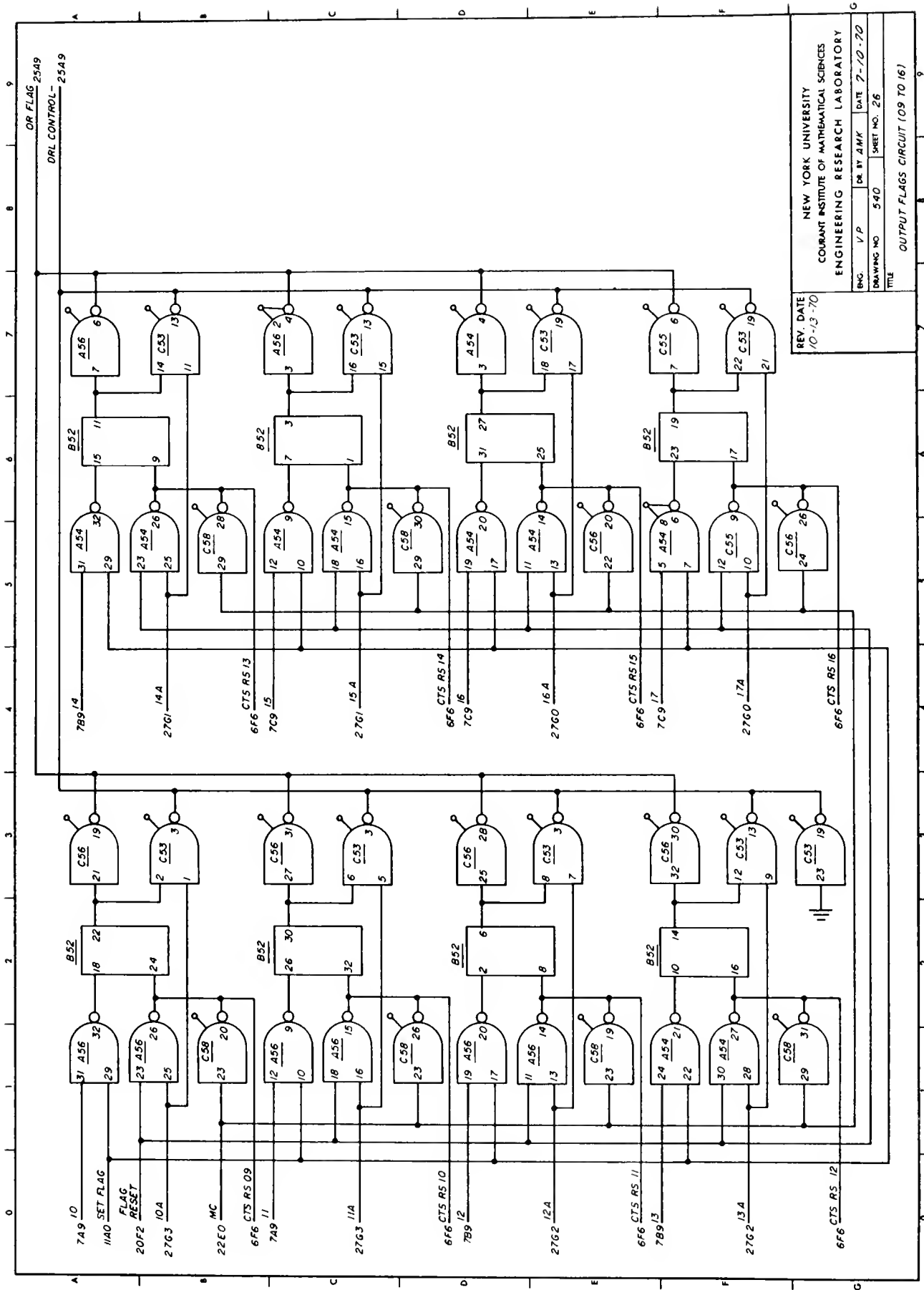




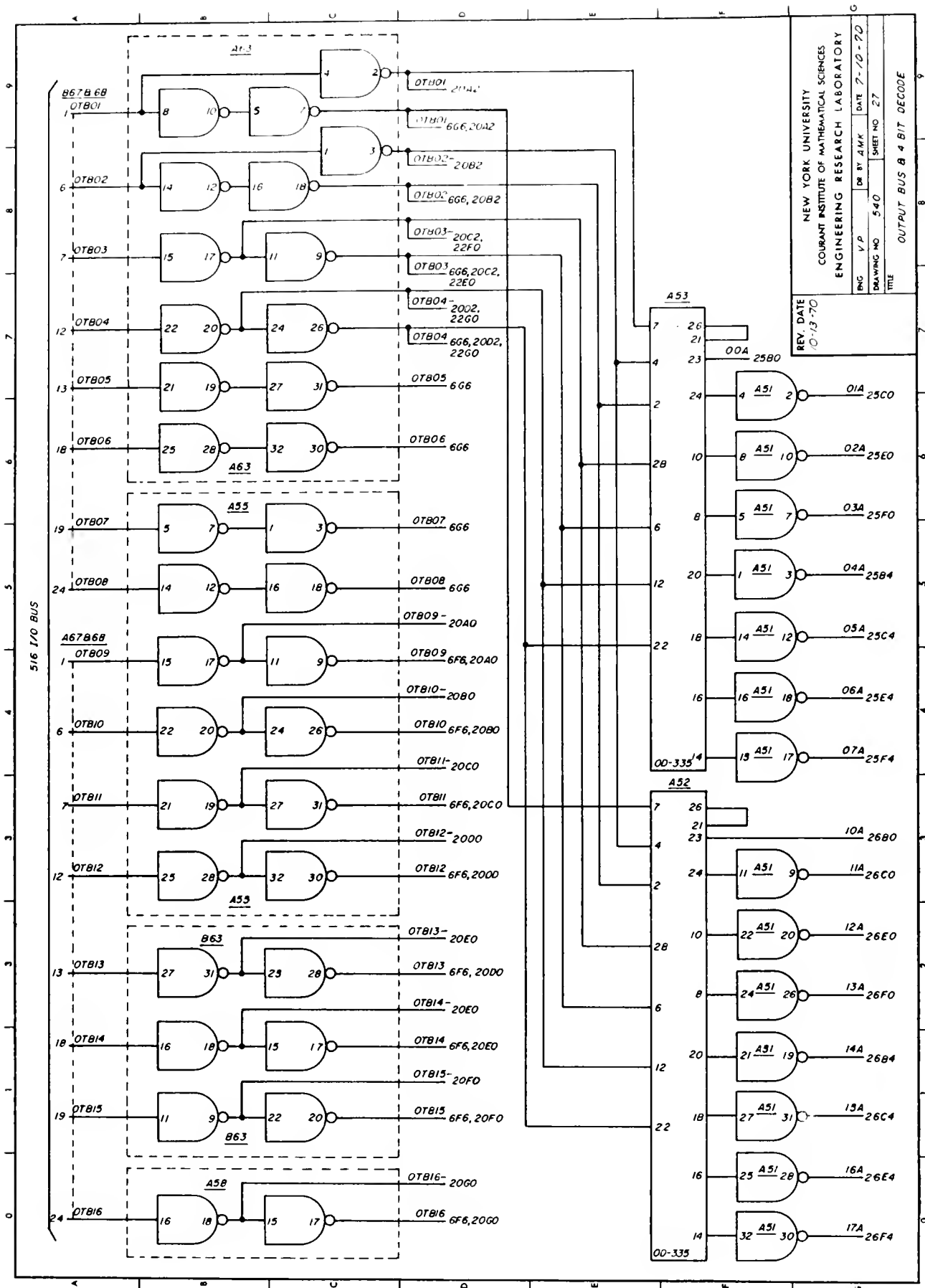
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ENG DRAWING NO	V.P. 54-0	DR BY ZMK	DATE 7-10-70
SHEET NO 25		TITLE OUTPUT FLAG CIRCUIT (OI TO OB)	







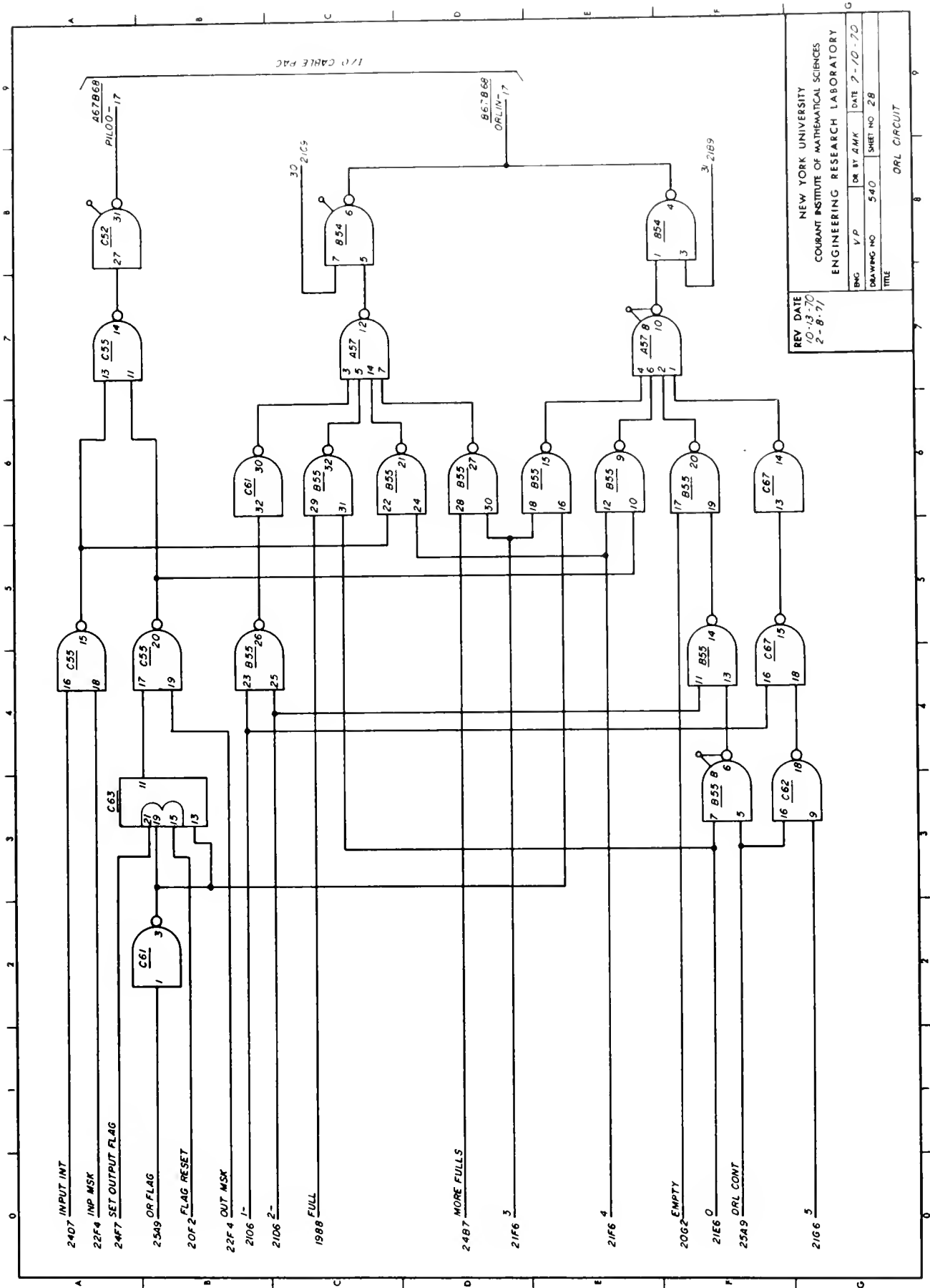




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SHEET NO. 27  
DATE 7-10-70  
TITLE OUTPUT BUS 4 BIT DECODE







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MULTIPLAR LAYOUT - RACK 1			

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HUB  
LL-PAC USED

BC-335	2
CC-152	3
CC-154	2
CM-022	2
DG-335	6
DG-336	6
DI-335	12
DJ-335	1
DL-335	13
FA-335	16
FF-335	2
MC-335	1
OD-335	4
PA-336	5
SM-330	6
TG-335	12
TOTAL	93

	A	B	C
1			
2			
3			
4			
5			
6			
7			
8			
2	DS 13-16 TO C58(PM)	FA-335	PA-336
3	DS 9-12 TO C48(PM)	FA-335	DL-335
4	DS 5-8 TO C38(PM)	TG-335	
5	DS 1-4 TO C28(PM)	DI-335	MC-335
6		TG-335	BC-335
7		OD-335	PA-336
8		OD-335	PA-336
3	STRIP SYNC CABLE PAC		
2	DG-335	DL-335	DL-335
3	DG-335	FA-335	FA-335
4	DG-335	TG-335	DI-335
5		DL-335	DL-335
6	TG-335	FA-335	FA-335
7	DL-335	DL-335	DL-335
8	DI-335	DI-335	DI-335
4	DG-336	DG-336	DG-336
2	DG-336	DG-336	DG-336
3			
4	SM-330	SM-330	SM-330
5			
6	SM-330	SM-330	SM-330
7	DJ-335		
8			BC-335
5	TG-335	DG-335	TG-335
2	OD-335	FF-335	CC-152
3	OD-335	FF-335	DG-335
4	DI-335	DI-335	PA-336
5	TG-335	DI-335	DI-335
6	DI-335	DI-335	CM022
7	DL-335	DL-335	DL-335
8	TG-335	TG-335	CM022
6	FA-335	FA-335	TG-335
2	FA-335	FA-335	DL-335
3	TG-335	TG-335	FA-335
4	FA-335	FA-335	FA-335
5	FA-335	FA-335	DG-335
6	CC-152	CC-152	PA-336
7	CC-154	CC-154	DI-335
8	I/O CABLE PAC (2)	I/O CABLE PAC (1)	DL-335

(VIEWED FROM LL-PAC SIDE)

RACK NO. 1





1	DS 788 CABLE	1	X6 532A	1	DS 586 CABLE
2	PM 788 TO C32	2	X6 532A	2	PM 586 TO C37
3	X7 532A	3	X8 532A	3	X5 532A
4	X7 532A	4	X8 532A	4	X5 532A
5	X1 532A	5	X4 532A	5	X3 532A
6	X1 532A	6	X4 532A	6	X3 532A
7	PM 182 TO C27	7	X2 532A	7	PM 384 TO C22
8	DS 182 CABLE	8	X2 532A	8	DS 384 CABLE
1	TG-335	1	DI-335	1	DI-335
2	DI-335	2	FA-335	2	CABLE TO C17
3	TG-335	3	TG-335	3	"Y"LC 542Y
4	FA-335	4	DI-335	4	"Y"LC 542Y
5	FA-335	5	FA-335	5	"Y"LC 542Y
6	DI-335	6	FA-335	6	"Y"LC 542Y
7	DI-335	7	FA-335	7	CABLE TO A17
8	DG-335	8	FA-335	8	CABLE(HUB A28)
1	TG-335	1	DI-335	1	DI-335
2	DI-335	2	FA-335	2	CABLE TO A12
3	TG-335	3	TG-335	3	"Y"LC 542Y
4	FA-335	4	DI-335	4	"Y"LC 542Y
5	FA-335	5	FA-335	5	FA-335
6	DI-335	6	FA-335	6	"Y"LC 542Y
7	DI-335	7	FA-335	7	CABLE TO C12
8	DG-335	8	FA-335	8	CABLE(HUB A26)
1	TG-335	1	DI-335	1	DI-335
2	DI-335	2	FA-335	2	CABLE TO C67
3	TG-335	3	TG-335	3	"Y"LC 542Y
4	FA-335	4	DI-335	4	"Y"LC 542Y
5	FA-335	5	FA-335	5	"Y"LC 542Y
6	DI-335	6	FA-335	6	"Y"LC 542Y
7	DI-335	7	FA-335	7	CABLE TO A67
8	DG-335	8	FA-335	8	CABLE(HUB A24)
1	TG-335	1	DI-335	1	DI-335
2	DI-335	2	FA-335	2	CABLE TO A62
3	TG-335	3	TG-335	3	"Y"LC 542Y
4	FA-335	4	DI-335	4	"Y"LC 542Y
5	FA-335	5	FA-335	5	"Y"LC 542Y
6	DI-335	6	FA-335	6	"Y"LC 542Y
7	DI-335	7	FA-335	7	DS 13814 CABLE TO C62
8	DG-335	8	FA-335	8	CABLE(HUB A22)
1	DS 15816 CABLE	1	X14 532A	1	DS 13814 CABLE
2	PM 15816 TO C52	2	X14 532A	2	PM 13814 TO C57
3	X15 532A	3	X16 532A	3	X13 532A
4	X15 532A	4	X16 532A	4	X13 532A
5	X9 532A	5	X12 532A	5	X11 532A
6	X9 532A	6	X12 532A	6	X11 532A
7	PM 9810 TO C47	7	X10 532A	7	PM(11B12) TO C42
8	DS 9810 CABLE	8	X10 532A	8	DS 11B12 CABLE

(VIEWED FROM 1L-PAC SIDE)

RACK NO 2

REV DATE  
10-13-70

NEW YORK UNIVERSITY  
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ENGINEERING RESEARCH LABORATORY

BNG V/P DATE 7-10-70  
DRAWING NO 540 SHEET NO 30

MULTIPLXR LAYOUT-RACK 2

PRE-MULTIPLXR  
1L-PAC USED  
PER 40 S  
DG-335 1  
DI-335 6  
FA-335 7  
TG-335 3  
"Y"LC(NYU 542Y) 4  
"X"LC(NYU 532A) 8



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NYO-1480-171  
Prestianni

c.2

AUTHOR

A 16 channel medium speed

TITLE

multiplex-r.

DATE DUE

BORROWER'S NAME

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Prestianni

c.2

AUTHOR

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TITLE

multi lex-r.

DATE DUE

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GERALD Feldman

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